

Hello DO!

An Overview of My Teaching and Research Activities

Mohamed Siala

<http://homepages.laas.fr/msiala>

17 avril 2019

INSA Toulouse, LAAS-CNRS

- This is not a technical presentation

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 - choose one paper to present ?
 - give a lecture about CP/SAT ?
 - talk about my research project ?
 - focus on potential collaborations ?
 - ...

Curriculum Vitae

- 2005 - 2007 : **Classes préparatoires en Mathématiques et Physiques**, IPEIS, Tunisia

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- 2007 - 2010 : **Engineering Diploma in Computer Science**, ENSI, Tunisia
- 2010 - 2012 : **Master's in Artificial Intelligence and Decision**, ENSI, Tunisia
- Decembre 2011, Mai 2015 : **PhD in Computer Science**, INSA Toulouse, LAAS-CNRS, ROC.

- 2011-2014 : **PhD. Candidate**, LAAS-CNRS, Toulouse, France,
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Supervisor : Barry O'Sullivan
Funding : Science Foundation Ireland (85%), and United
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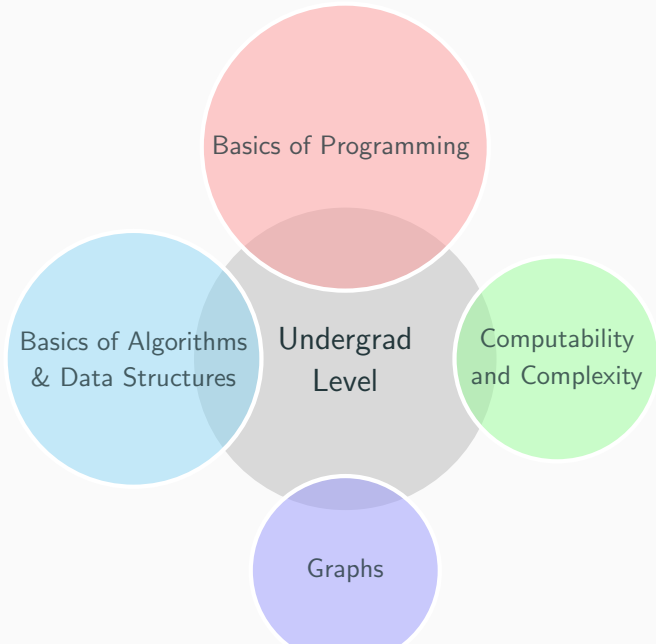
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- Since 08/2018 : **Associate Professor**, INSA Toulouse, LAAS-CNRS

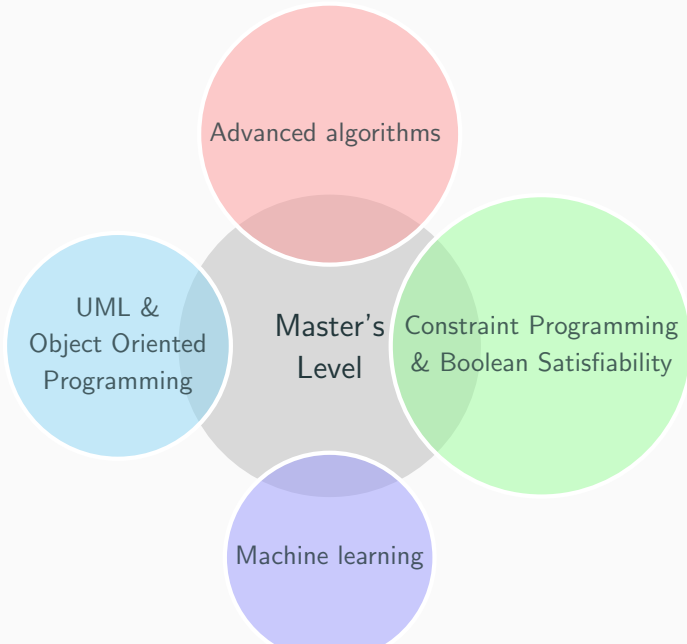
An overview of my Teaching activities

- INSA Toulouse : Engineering School (5 years),
- Department of Electrical and Computer Engineering
- Involved in teaching courses from 1st year to 5th year

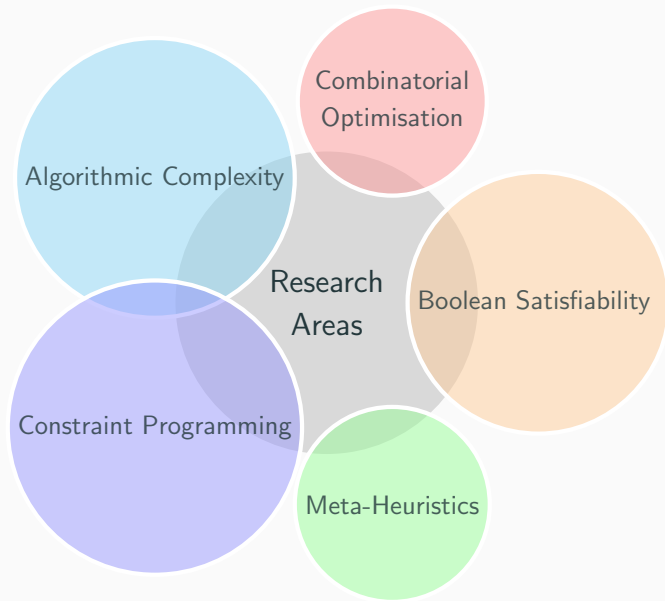
Undergrad Level (1st, 2nd, 3rd year)



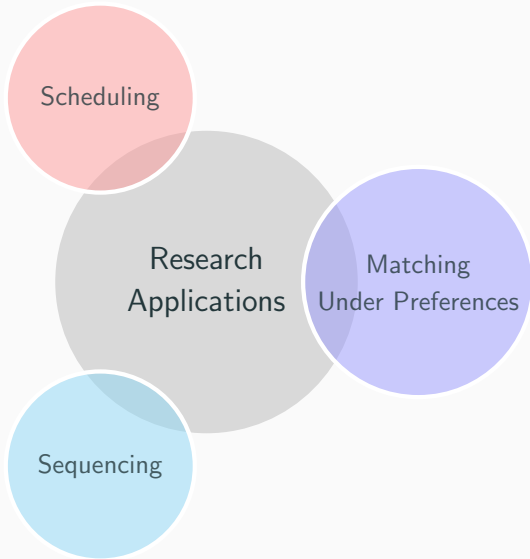
Master's Level (4th and 5th year)



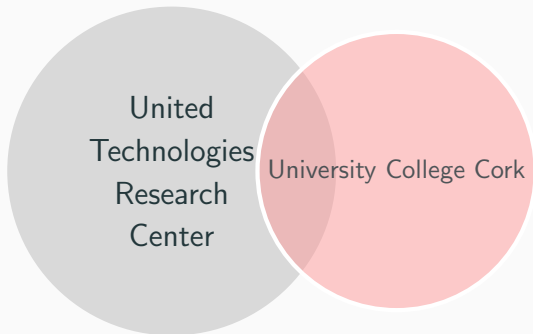
An overview of past research



An overview of past research applications



An overview of past industrial collaborations : UCC-UTRC



Begum Genc : Supervision with Barry O'Sullivan.

An Approach to Robustness in Stable Marriage and Stable Roommates Problems

PhD defence : May 3rd 2019

An Introduction to Constraint Programming

- Hard computational problems
- Combinatorial problems

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Solving Combinatorial Problems

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- Metaheuristics

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 - ⊖ Not flexible
- Specific modeling language (MIP/SAT)
 - ⊖ Modelisation matters
- Metaheuristics
 - ⊕ fast approximation
 - ⊖ Do not guarantee optimal solution
 - ⊖ Decision problems?

Constraint programming

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- The user states the problem, the computer finds a solution

Constraint programming

- Declarative approach
- The user states the problem, the computer finds a solution
- Clear separation between modeling and solving.

Constraint Network

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Constraint Satisfaction Problem

Is there an assignment to the variables that satisfy all the constraints?

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Constraint Optimisation Problem

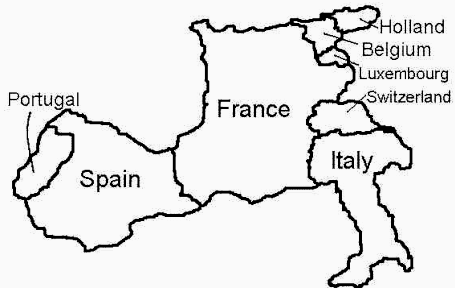
Is there an assignment to the variables that satisfy all the constraints *and optimise an objective function*?

Important note

Observe that no restriction on the constraints is made. A constraint could be

- Logical formula
- Arithmetic
- Linear
- Non Linear
- Tabular
- A conjunction of other constraints

Example : Map coloring



Example : Map coloring

Model

- Variables :
Portugal, Spain, France, Italy, Switzerland, Luxembourg, Belgium, Holland
- Domains : {*Black, Blue, Green, Red, White, Yellow*}
- Constraints :
 1. *Portugal* \neq *Spain*
 2. *Spain* \neq *France*
 3. *France* \neq *Italy*
 4. *France* \neq *Switzerland*
 5. *France* \neq *Luxembourg*
 6. *France* \neq *Belgium*
 7. *Belgium* \neq *Holland*
 8. *Belgium* \neq *Luxembourg*
 9. *Switzerland* \neq *Luxembourg*
 10. *Switzerland* \neq *Italy*

Philosophy of Constraint Programming

- In order to find solutions, eliminate as much possible non-solutions
[filtering]

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Example : Sudoku

8			4		6			7
						4		
	1					6	5	
5		9		3		7	8	
				7				
	4	8		2		1		3
	5	2					9	
		1						
3			9		2			5

7	2	5
3	1	8
?	4	9

7	2	5
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6	4	9

Filtering

7	?	5	
3	1	?	
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Filtering

7	{2, 6, 8}	5	
3	1	{2, 6, 8}	
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- Decisions : Domain operations ($x = v$, $x \leq v$, ...)

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- We have to come up with a filtering algorithm for each sub-problem
- We usually know how to solve such sub-problem individually. However, we need to reason about them from a filtering point of view.
- Excellent strategy if the sub-problem is recurrent! \implies we can use the same constraint in many problems

**Example of a Research Paper :
Rotation-Based Formulation for
Stable Matching
Mohamed Siala and Barry
O'Sullivan, CP-'17, August 2017,
Melbourne, Australia**

Matching Under Preferences

Matching Under Preferences



- Assign residents to hospitals
- Every resident has a personnel preference over hospitals
- Each hospital has a preference list over residents

Matching Under Preferences



- Assign students to universities
- Every student has a personal preference over universities
- Each university has a ranking list over students

Many-to-Many Stable Matching

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Example of Stability

w_1	$f_1 f_2 f_3$
w_2	$f_2 f_3 f_1$
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w_2	$f_2 f_3 f_1$	f_2	$w_1 w_2 w_3$
w_3	$f_3 f_1 f_2$	f_3	$w_3 w_1 w_2$

- $M_1 = \{\langle w_1, f_2 \rangle, \langle w_2, f_3 \rangle, \langle w_3, f_1 \rangle\}$ is not stable because $\langle w_3, f_3 \rangle$ is blocking M_1

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w_1	$f_1 f_2 f_3$	f_1	$w_3 w_2 w_1$
w_2	$f_2 f_3 f_1$	f_2	$w_1 w_2 w_3$
w_3	$f_3 f_1 f_2$	f_3	$w_3 w_1 w_2$

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- $M_2 = \{\langle w_1, f_2 \rangle, \langle w_2, f_1 \rangle, \langle w_3, f_3 \rangle\}$ is stable

Hard Variants

- Many-2-Many stable matching is solvable in $O(L)$ time
- However, the problem becomes NP-Hard with side constraints

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Sex-Equal & Balanced Stable Matching

- Let M be a stable matching
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- Balanced stable matching : find a stable matching M with the minimum value of $\max\{C_M^w, C_M^f\}$
- Sex-Equal Stable matching : find a stable matching M with the minimum value of $|C_M^m - C_M^w|$

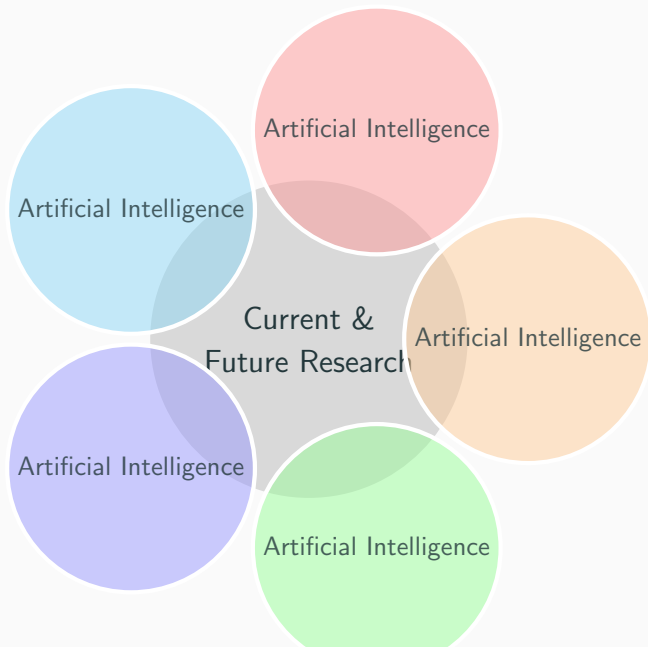
Stable Matching via CP

- Stability as a constraint
- Reformulation as a set of clauses
- We showed that arc consistency can be maintained in $O(L^2)$ worst case time complexity

Future Research and Potential Collaborations in DO

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 - MAC : Applications of CP in continuous optimisation ?

Thank You !

Awards

- European Association for Artificial Intelligence Distinguished Dissertation Award 2015, honourable mention
- XCSP3 Solver Competition 2017 : First place in the optimisation track
- Honorable mention for the paper "An Optimal Arc Consistency Algorithm for a Chain of Atmost Constraints with Cardinality", CP2012 Conference