



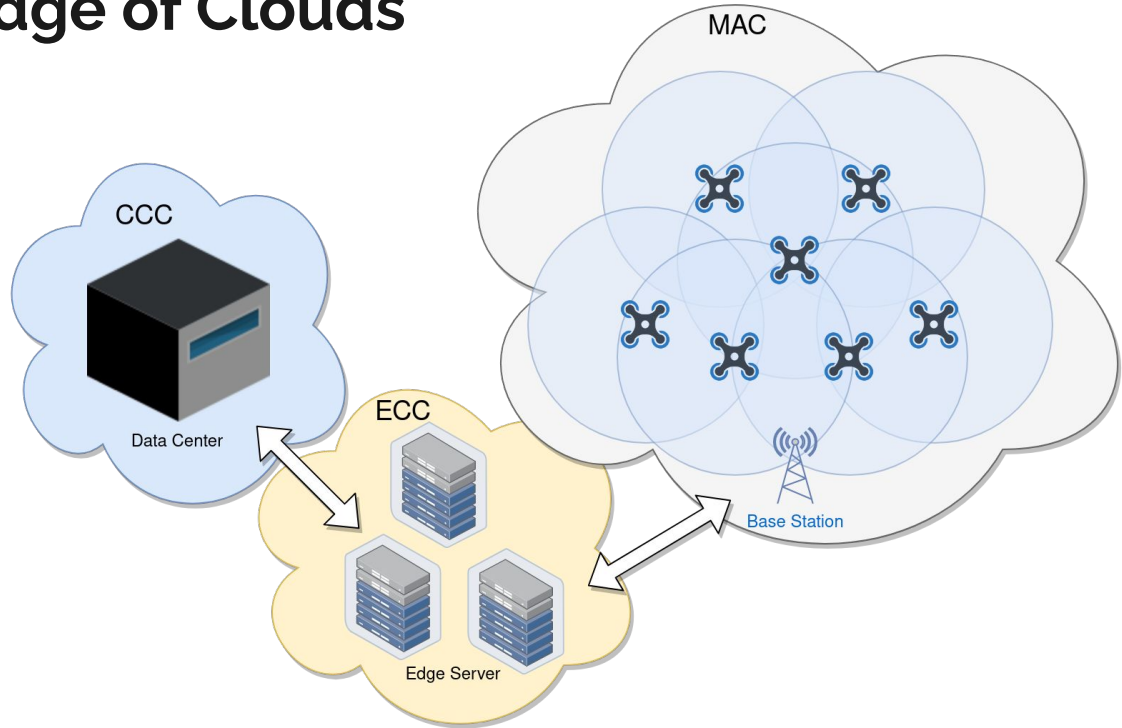
# Assessing Traffic Flow and Using Mobility for Distributed Applications Performance Enhancement.

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# Computing at the Edge of Clouds

- Applications physically far from cloud infrastructures can benefit from closer edge servers.
- Data can be pre-processed at the edge and only processed data sent to central clouds
- Computational power at the edge can be leveraged for constrained nodes
- Improved security and privacy
- Mobile ad-hoc computing





# Motivation

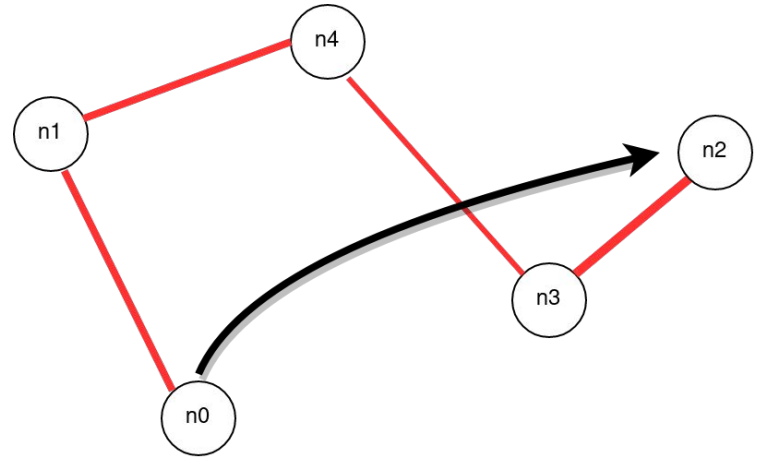
- Applications such as remote sensing and earth observation with swarms of UAVs or Satellites
- Connectivity issues of Mobile ad-hoc multihop systems
  - Unstable connections
  - Lower bandwidth
  - Channel sharing
  - MAC overhead
- Nodes are both source, destination and routers
- How the position of the nodes in the network affect the performance of the applications?



# Working on the Application Layer

When we think about network congestion, it is natural to think how a routing protocol could solve that.

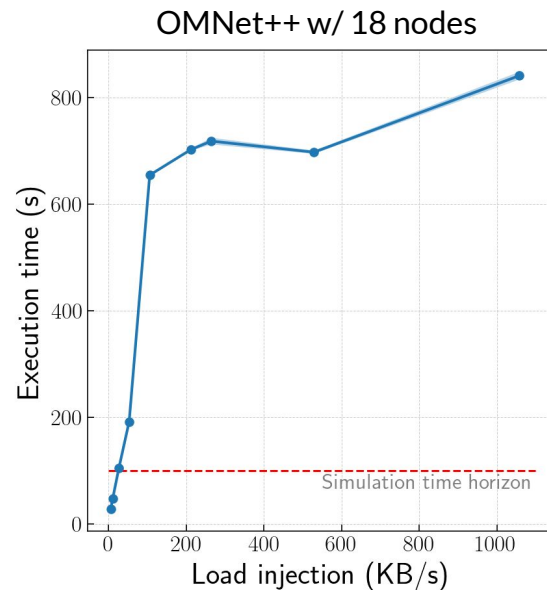
- Fragmentation of routing protocols
- Ultra specialization of protocols



# Existing Tools don't Scale Well

Fine grained discrete simulations usually don't scale because they are resource hungry

- Too much memory and too much processing power to simulate each message, or packet.
- One alternative solution is to use macroscopic fluid models.





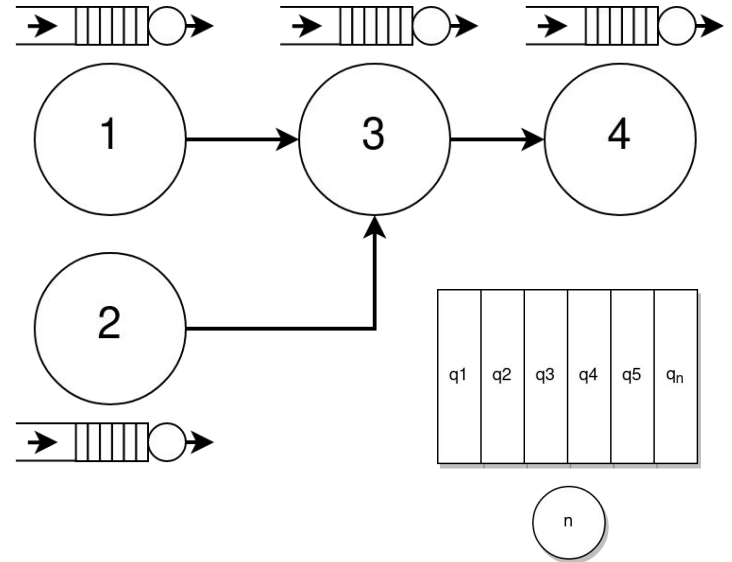
# Network Model and Results

- Network Model Introduction
  - Throughput
  - Latency
- Topologies and Mobility
- Results
- Can we enhance performance by using mobility?

# Traffic Flow Model

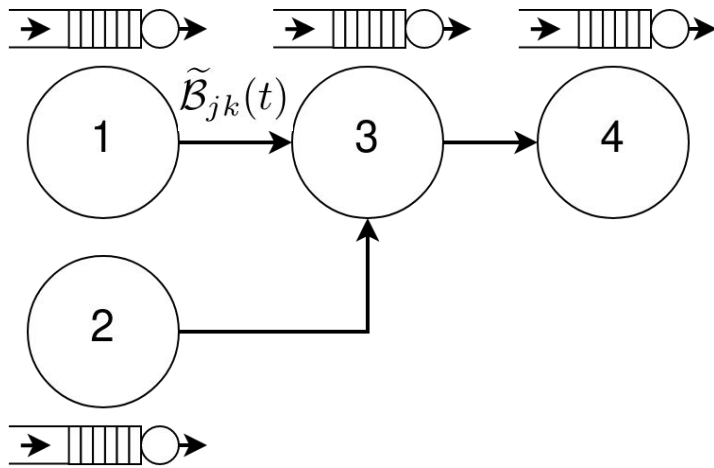
One ad hoc network topology can be modeled as a multi-queue system:

- For each simulation step, part of each queue leaves the node according to the amount of available bandwidth.
- The forwarded data is added to next hop's queue and moves towards the final destination.





# Traffic Flow Model



Departure 
$$d_j^k(t) = \begin{cases} -\frac{Q_j^k(t)}{Q_j(t)} \tilde{B}_{jk}(t), & \text{if } Q_j^k(t) > 0 \\ 0, & \text{otherwise} \end{cases}$$

Arrival 
$$a_{e=k}^j(t) = \begin{cases} \sum_{\substack{i=0 \\ i \neq j}}^{n-1} \frac{Q_i^j(t)}{Q_i(t)} \tilde{B}_{ij}(t), & \text{if } Q_i^j(t) > 0 \\ 0, & \text{otherwise} \end{cases}$$

$$Q_j^k(t) = -\frac{Q_j^k(t)}{Q_j(t)} \tilde{B}_{jk}(t) \mathbb{1}_{Q_j^k(t) > 0} \mathbb{1}_{j \neq k} + \sum_{\substack{i=0 \\ i \neq j}}^{n-1} \frac{Q_i^k(t)}{Q_i(t)} \tilde{B}_{ij}(t) \mathbb{1}_{Q_i^j(t) > 0} \mathbb{1}_{\mathcal{R}_{i,j}(t)=j}$$



# Using Average Age

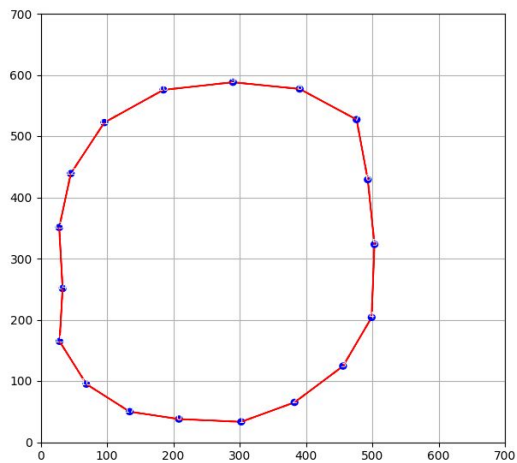
As a way to model Latency

In order to model the latency of each message, conservation of momentum is used with messages' age being the conserved quantity.

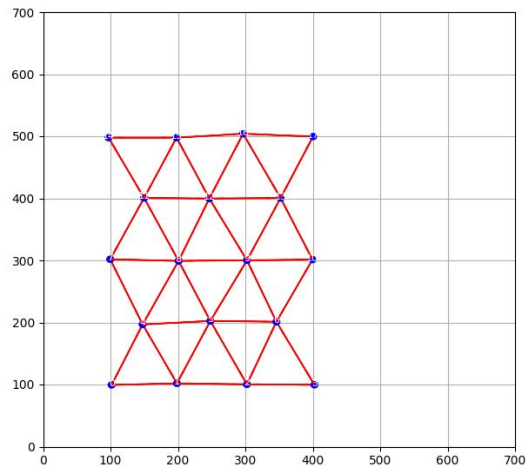
$$\begin{aligned} & [(\mathcal{A}_j^k(t - \delta t) + \delta t)(\mathcal{Q}_j^k(t - \delta t) - d_j^k(t - \delta t)) + \\ & + \sum_{\substack{i=0 \\ i \neq j \\ \mathcal{H}_i^k(t)=j}}^{n-1} (\mathcal{A}_i^k(t - \delta t) + \delta t)d_i^k(t - \delta t)] - \mathcal{A}_j^k(t)\mathcal{Q}_j^k(t) = 0 \end{aligned}$$



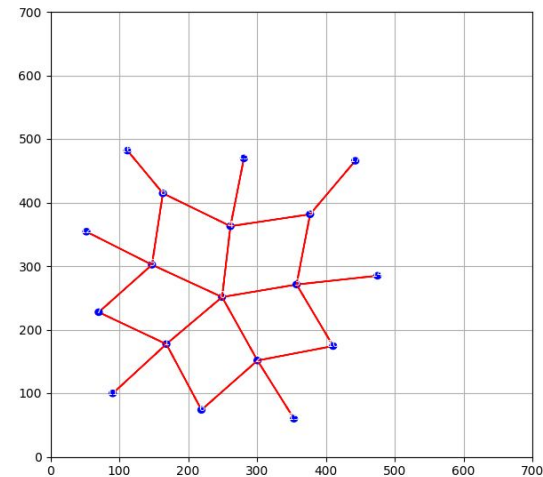
# Simulated Topologies



Ring



Crystal

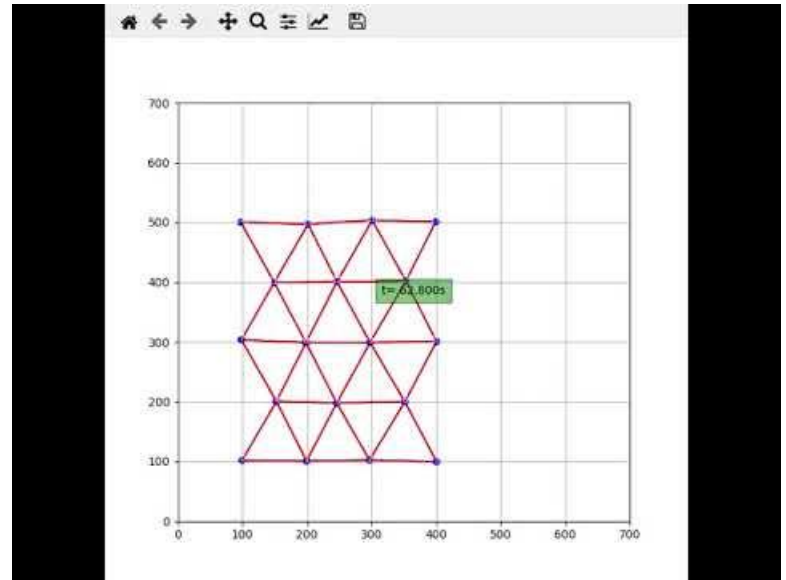


Star

# Simulation Results

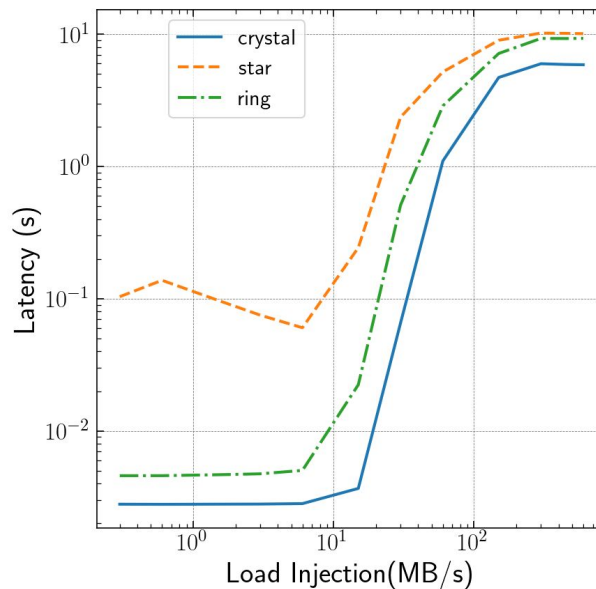
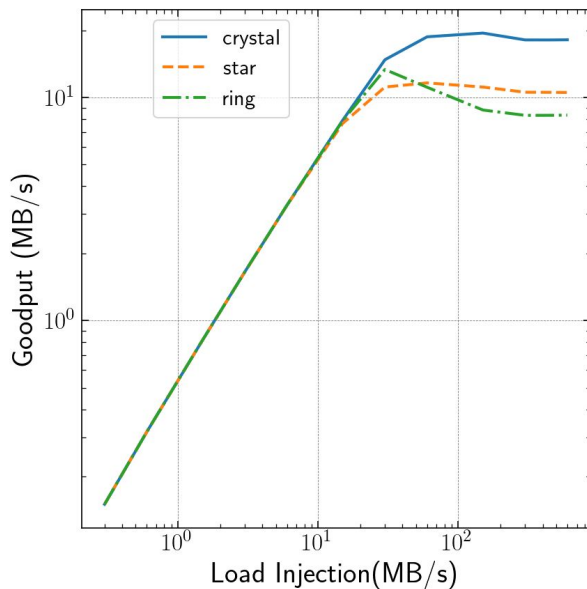
## Data Injection profile

- Round Robin data injection
- $f$  = variable, data = 3MB
- One to all sharing of equal chunks
- $t_h = 100s$



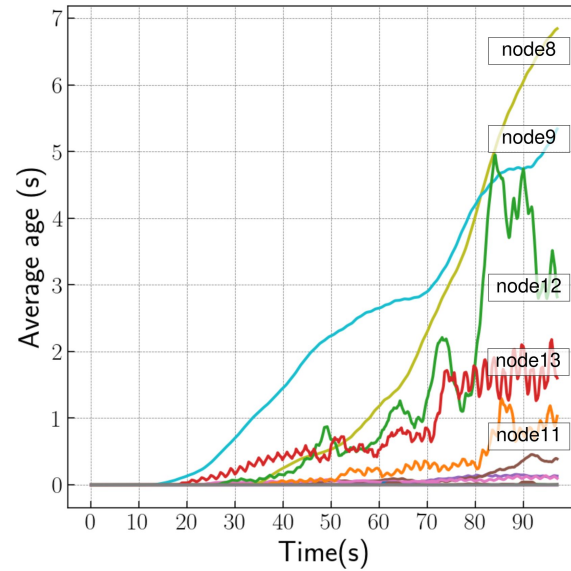
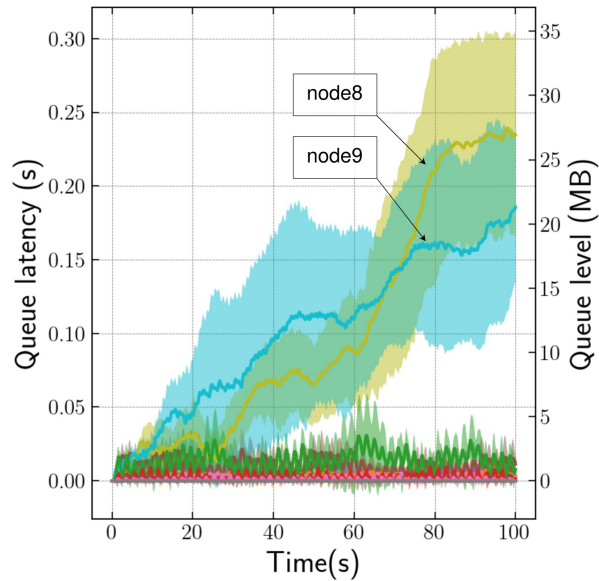
# Choosing the Best Topology for the Traffic Pattern

## Network Saturation

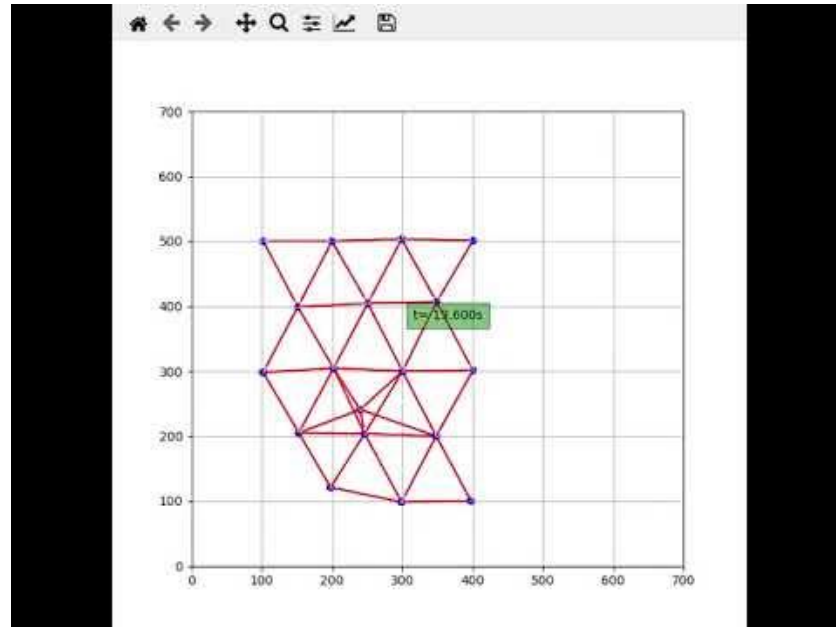


# Congestion and Latency

Round Robin 30MB/s



# Can we Support Congested Areas?

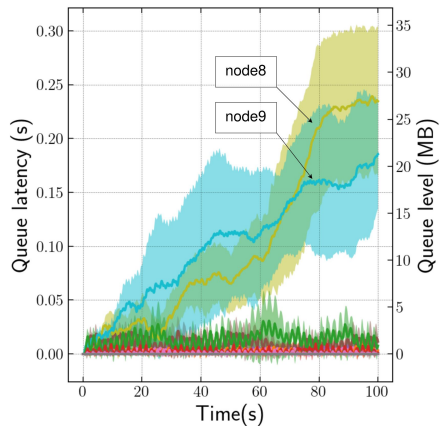




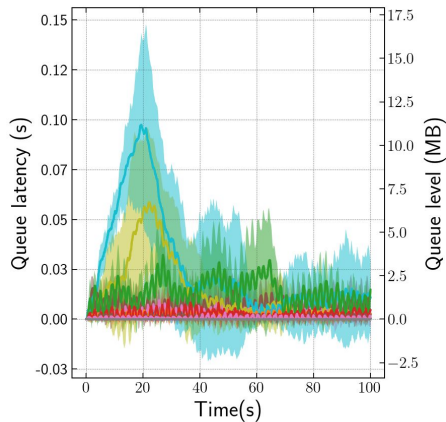
# Congestion and Latency

## Round Robin 30MB/s

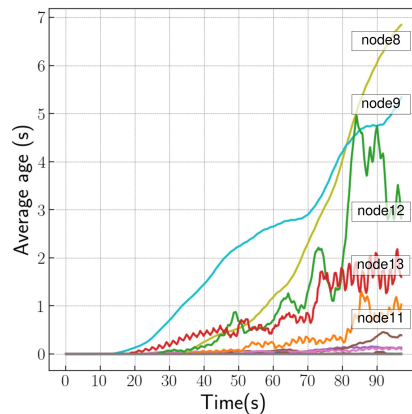
Without Support



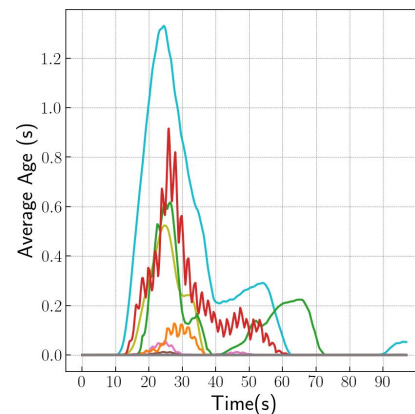
With Support



Without Support

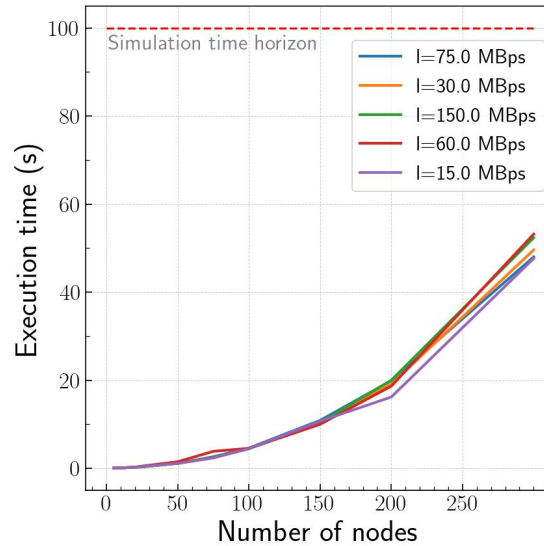


With Support



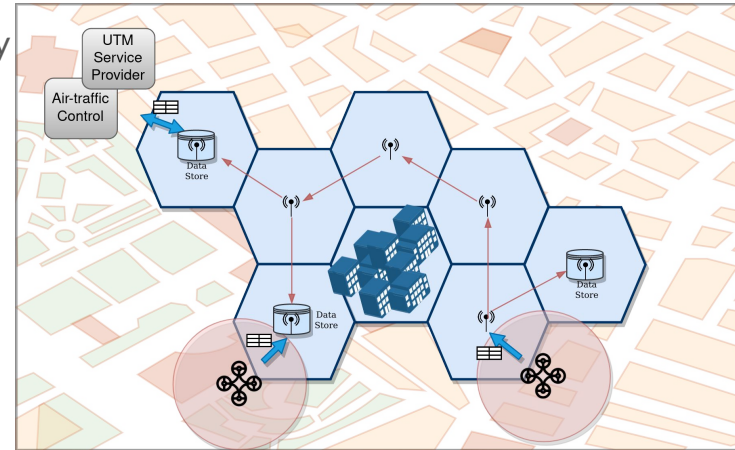


# Scalability Assessment



# Towards a Middleware

- Establish heuristics and control strategies to use the mobility in our favor, to reduce congestion and latency and increase system performance.
- Assess how to use those strategies to find better placement for replicas according to mobile node density across the topology.
  - Use these placement strategies to support other ongoing research about the use of distributed stores to support UAV position tracking for UTM





# Conclusion

- We can model network traffic flow in Mobile Ad-hoc Computing with a lightweight model and measure the most common metrics in distributed applications such as throughput and latency.
- We can observe the influence of different topologies in application performance and choose the one that yields higher throughput and lower latency.
- We can also influence the and potentially lower the formation of congestion in the topologies by using the mobility of the nodes during runtime.



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