

KRR2: First-order logic

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- 1 First-order logic, why?
- 2 Syntax and semantics

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Pros and cons of propositional logic

Pros

- Propositional logic is **declarative**: pieces of syntax correspond to facts
- Propositional logic allows partial/disjunctive/negated information
(unlike most data structures and databases)
- Propositional logic is **compositional**: meaning of $a \wedge b$ is derived from meaning of a and of b
- Meaning in propositional logic is **context-independent**
(unlike natural language, where meaning depends on context)

Cons

Propositional logic has very limited expressive power (unlike natural language)

Expressiveness of propositional logic

Example

I want to declare:

Every friend of my sisters has a blue car

In the natural language, I need one simple sentence. What about in propositional logic?

I need symbols!! Lots of them!! Because I have a big family and my sisters are very friendly.

- Sister_1_Friend_1_Has_Blue_Car
- Sister_1_Friend_2_Has_Blue_Car
- ...
- Sister_4_Friend_23_Has_Blue_Car

First-order logic

Whereas propositional logic assumes world contains *facts*, first-order logic (like natural language) assumes the world contains

- **Objects:** people, houses, numbers, theories, colors, cricket games, centuries . . . and me, and cars!!
- **Relations:** red, round, bogus, prime, multistoried . . . , brother of, bigger than, inside, part of, has color, occurred after, owns, comes between, . . . and blue!!
- **Functions:** father of, third inning of, one more than, end of . . . and friend of, sister of!!

Outline

- 1 First-order logic, why?
- 2 Syntax and semantics**

Syntax of First-Order Logic

Basic elements

- 1 **Constants:** *KingJohn, 2, ANU, Yannick* ...
- 2 **Predicate:** *Sister, >* ...
- 3 **Functions:** *Sqrt, FriendOf* ...
- 4 **Variables:** *x, y, a, b* ...
- 5 **Connectives:** $\vee, \wedge, \neg, \Rightarrow, \Leftrightarrow$
- 6 **Equality:** $=$
- 7 **Quantifiers:** $\forall \exists$

Syntax of First-Order Logic

Term

A **term** represents an object in FOL. Its syntax is:

- a constant, or
- a variable, or
- a function of terms $function(term_1, \dots, term_n)$

Atomic sentence

An **atomic sentence** represents an elementary relation between terms. Its syntax is:

- a predicate $predicate(term_1, \dots, term_n)$
- an equality of terms $term_1 = term_2$

Example

$Brother(KingJohn, RichardTheLionheart)$
 $> (Length(LeftLegOf(Richard)), Length(LeftLegOf(KingJohn)))$
 $carOf(friendOf(oneSisterOf(Yannick))) = colorOf(Ocean)$

Syntax of First-Order Logic

Complex sentences

Complex sentences are made from atomic sentences using connectives

$$\neg S, \quad S_1 \wedge S_2, \quad S_1 \vee S_2, \quad S_1 \Rightarrow S_2, \quad S_1 \Leftrightarrow S_2$$

Example

$Sibling(KingJohn, Richard) \Rightarrow Sibling(Richard, KingJohn)$

$>(1, 2) \vee \leq(1, 2)$

$>(1, 2) \wedge \neg >(1, 2)$

$Sister(Marie, Yannick) \Rightarrow CarColor(FriendOf(Marie), blue)$

Truth in first-order logic

Semantics

Sentences are true with respect to a **model** and an **interpretation**.

Model

Model contains objects (**domain elements**) and relations among them.

Interpretation

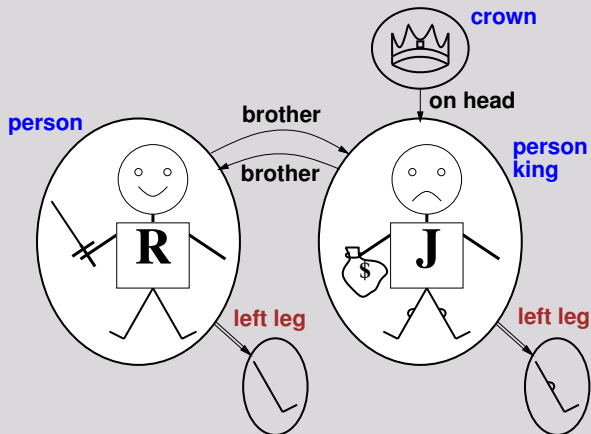
Interpretation specifies referents for

- *constant symbols* → objects
- *predicate symbols* → relations
- *function symbols* → functional relations

An atomic sentence $\text{predicate}(term_1, \dots, term_n)$ is true iff the objects referred to by $term_1, \dots, term_n$ are in the relation referred to by *predicate*.

Models for FOL: Example

Example



Models for FOL: Example

Example

Consider the interpretation in which

- *Richard* → Richard the lionheart
- *John* → the evil King John
- *Brother* → the brotherhood relation

Under this interpretation, *Brother(Richard, John)* is true just in case Richard the Lionheart and the evil King John are in the brotherhood relation in the model

Example

Consider this new interpretation on the model about Yannick and his sisters in which

- *Richard* → Yannick
- *John* → Marie
- *Brother* → the brother-sister relation

Under this interpretation, *Brother(Richard, John)* is true just in case Yannick and Marie are in the brother-sister relation in the model.

This interpretation is ok but the symbols are not very well-chosen!

Universal quantification

Universal quantification

Symbol: \forall

Syntax: \forall $\langle variables \rangle$ $\langle sentence \rangle$

Semantics: $\forall x$ P is true in a model m iff P is true with x being **each possible object** in the model

Example

Quantified sentence: $\forall x$ $Sister(x, Yannick) \Rightarrow ColorCar(FriendOf(x), blue)$

Roughly speaking, it is equivalent to:

$Sister(Marie, Yannick) \Rightarrow ColorCar(FriendOf(Marie), Blue)$

$\wedge Sister(Claire, Yannick) \Rightarrow ColorCar(FriendOf(Claire), Blue)$

$\wedge Sister(KingJohn, Yannick) \Rightarrow ColorCar(FriendOf(KingJohn), Blue)$

$\wedge Sister(Blue, Yannick) \Rightarrow ColorCar(FriendOf(Blue), Blue)$

$\wedge Sister(Yannick, Yannick) \Rightarrow ColorCar(FriendOf(Yannick), Blue) \wedge \dots$

A common mistake to avoid

Main connective with \forall

Typically, \Rightarrow is the main connective with \forall

Common mistake: using \wedge as the main connective with \forall

Example

$$\forall x \text{ Sister}(x, \text{Yannick}) \wedge \text{ColorCar}(\text{FriendOf}(x), \text{Blue})$$

means “Everyone is the sister of Yannick and every friend of everyone has a blue car”

Existential quantification

Existential quantification

Symbol: \exists

Syntax: $\exists < \text{variables} > < \text{sentence} >$

Semantics: $\exists x$ P is true in the model m iff P is true with x being some object in the model.

Example

Quantified sentence: $\exists x \text{ Sister}(x, \text{Yannick}) \wedge \text{ColorCar}(\text{FriendOf}(x), \text{blue})$

Roughly speaking, it is equivalent to:

$\text{Sister}(\text{Marie}, \text{Yannick}) \wedge \text{ColorCar}(\text{FriendOf}(\text{Marie}), \text{Blue})$

$\vee \text{Sister}(\text{Claire}, \text{Yannick}) \wedge \text{ColorCar}(\text{FriendOf}(\text{Claire}), \text{Blue})$

$\vee \text{Sister}(\text{KingJohn}, \text{Yannick}) \wedge \text{ColorCar}(\text{FriendOf}(\text{KingJohn}), \text{Blue})$

$\vee \text{Sister}(\text{Blue}, \text{Yannick}) \wedge \text{ColorCar}(\text{FriendOf}(\text{Blue}), \text{Blue})$

$\vee \text{Sister}(\text{Yannick}, \text{Yannick}) \wedge \text{ColorCar}(\text{FriendOf}(\text{Yannick}), \text{Blue}) \vee \dots$

Another common mistake to avoid

Main connective with \exists

Typically, \wedge is the main connective with \exists

Common mistake: using \Rightarrow as the main connective with \exists

Example

$$\exists x \text{ Sister}(x, \text{Yannick}) \Rightarrow \text{ColorCar}(\text{FriendOf}(x), \text{Blue})$$

This sentence is true if you find someone who is not my sister!
It does not matter if this person has a blue car or not.

Properties of quantifiers

Properties

- $\forall x \forall y$ is the same as $\forall y \forall x$
- $\exists x \exists y$ is the same as $\exists y \exists x$
- $\exists x \forall y$ is **not** the same as $\forall y \exists x$
 - $\exists x \forall y \text{ Loves}(x, y)$
 - “There is a person who loves everyone in the world”
 - $\forall y \exists x \text{ Loves}(x, y)$
 - “Everyone in the world is loved by at least one person”

Quantifier duality

Each quantifier can be expressed using the other

- $\forall x \text{ Likes}(x, \text{IceCream})$ is the same as $\neg \exists x \neg \text{Likes}(x, \text{IceCream})$
- $\exists x \text{ Likes}(x, \text{Broccoli})$ is the same as $\neg \forall x \neg \text{Likes}(x, \text{Broccoli})$

Fun with sentences

Example

- Brothers are siblings

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$$\forall x, y \text{ Brother}(x, y) \Rightarrow \text{Sibling}(x, y)$$

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- “One’s mother is one’s female parent”

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- “One’s mother is one’s female parent”

$$\forall x, y \text{ Mother}(x, y) \Leftrightarrow (\text{Female}(x) \wedge \text{Parent}(x, y))$$

Example

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- “Sibling” is symmetric

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- “One’s mother is one’s female parent”

$$\forall x, y \text{ Mother}(x, y) \Leftrightarrow (\text{Female}(x) \wedge \text{Parent}(x, y))$$

- “A first cousin is a child of a parent’s sibling”

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- “Sibling” is symmetric

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- “One’s mother is one’s female parent”

$$\forall x, y \text{ Mother}(x, y) \Leftrightarrow (\text{Female}(x) \wedge \text{Parent}(x, y))$$

- “A first cousin is a child of a parent’s sibling”

$$\forall x, y \text{ FirstCousin}(x, y) \Leftrightarrow \exists p, ps \text{ Parent}(p, x) \wedge \text{Sibling}(ps, p) \wedge \text{Parent}(ps, y)$$

Equality

Equality

$term_1 = term_2$ is true under a given interpretation if and only if $term_1$ and $term_2$ refer to the same object.

Example

- $1 = 2$ is satisfiable (if the symbols 1 and 2 refer to the same object in the interpretation)
- $2 = 2$ is valid

Example

Definition of *Sibling* thanks to *Parent*:

$$\forall x, y \text{ Sibling}(x, y) \Leftrightarrow (\neg(x = y) \wedge \exists m, f \neg(m = f) \wedge$$

$$\text{Parent}(m, x) \wedge \text{Parent}(f, x) \wedge \text{Parent}(m, y) \wedge \text{Parent}(f, y))$$

Summary

- Knowledge representation language:
 - declarative, compositional, expressive, context-independent, unambiguous
- Model: set of objects, functions and their relation
- Knowledge-base in first-order logic
 - careful process
 - 1 analyzing the **domain** (objects, functions, relations),
 - 2 choosing a **vocabulary** (interpretation)
 - 3 encoding the **axioms** (what is known in *KB*) to support the desired **inferences**