

COURSE ON LMI OPTIMIZATION
WITH APPLICATIONS IN CONTROL
PART I.2

GEOMETRY OF LMI SETS

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Macchina Comica (1928)
František Kupka (1871-1957)

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Geometry of LMI sets

Given symmetric matrices F_i we want to characterize the shape in \mathbb{R}^n of the LMI set

$$\mathcal{F} = \{x \in \mathbb{R}^n : F(x) = F_0 + \sum_{i=1}^n x_i F_i \succeq 0\}$$

Matrix $F(x)$ is PSD iff its **diagonal minors** $f_i(x)$ are nonnegative

Diagonal minors are multivariate **polynomials** of indeterminates x_i

So the LMI set can be described as

$$\mathcal{F} = \{x \in \mathbb{R}^n : f_i(x) \geq 0, i = 1, \dots, n\}$$

which is a **semialgebraic** set

Moreover, it is a **convex** set

Example of 2D LMI feasible set

$$F(x) = \begin{bmatrix} 1 - x_1 & x_1 + x_2 & x_1 \\ x_1 + x_2 & 2 - x_2 & 0 \\ x_1 & 0 & 1 + x_2 \end{bmatrix} \succeq 0$$

Feasible iff all principal minors nonnegative

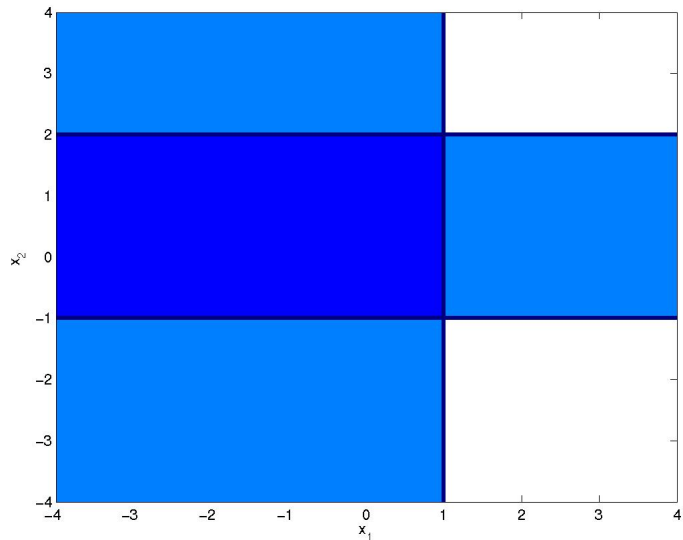
System of **polynomial inequalities** $f_i(x) \geq 0$

1st order minors

$$f_1(x) = 1 - x_1 \geq 0$$

$$f_2(x) = 2 - x_2 \geq 0$$

$$f_3(x) = 1 + x_2 \geq 0$$

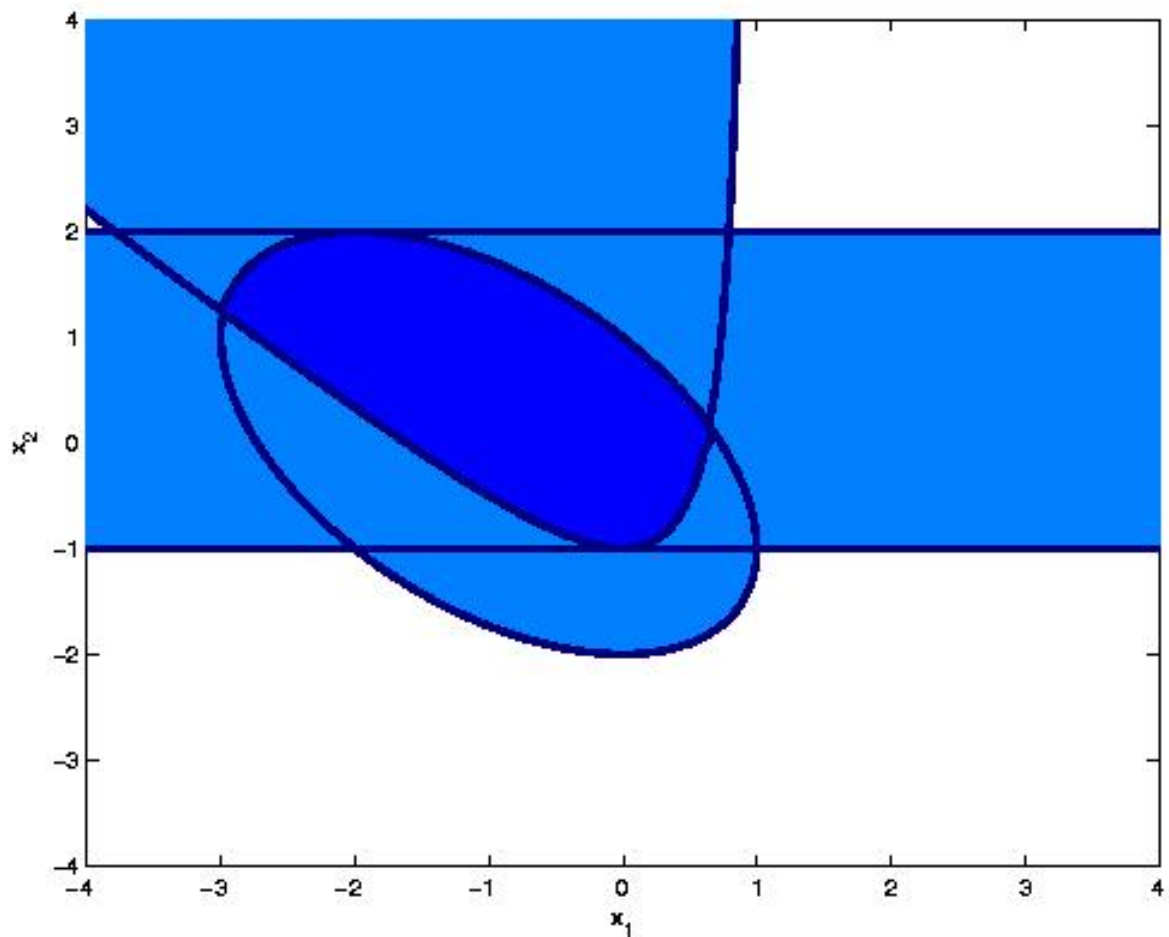


2nd order minors

$$f_4(x) = (1 - x_1)(2 - x_2) - (x_1 + x_2)^2 \geq 0$$

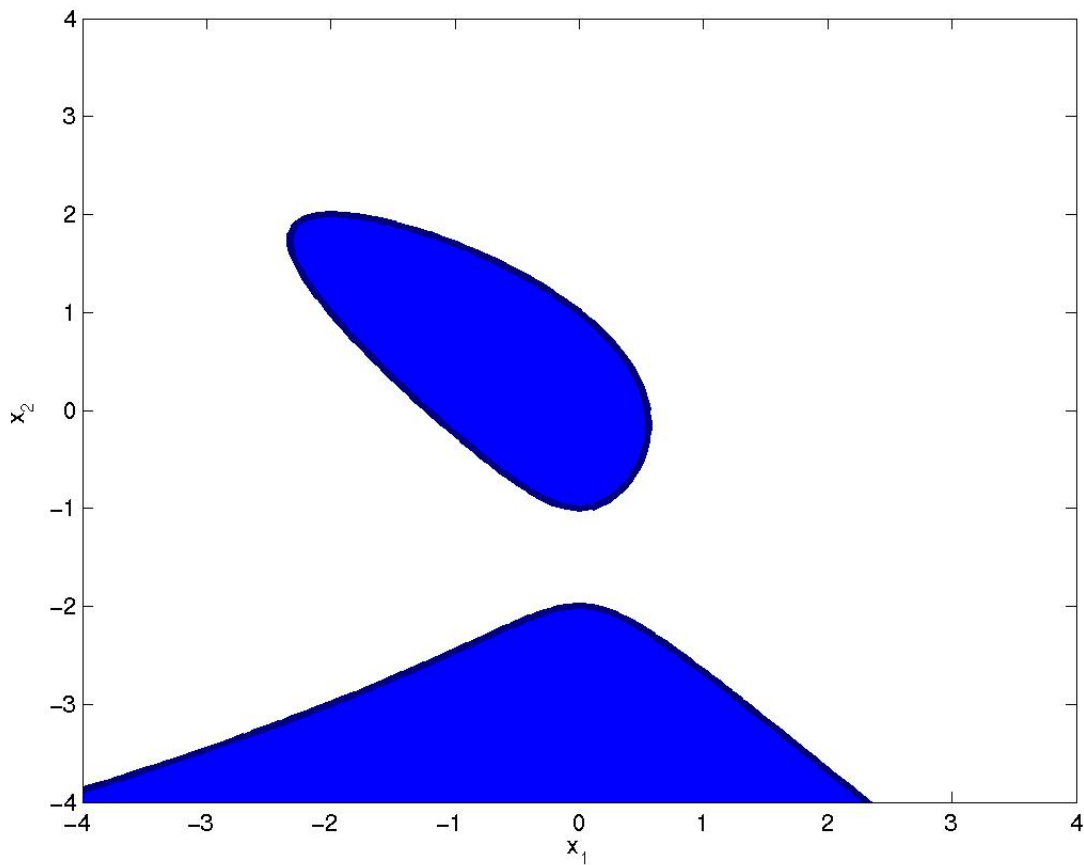
$$f_5(x) = (1 - x_1)(1 + x_2) - x_1^2 \geq 0$$

$$f_6(x) = (2 - x_2)(1 + x_2) \geq 0$$

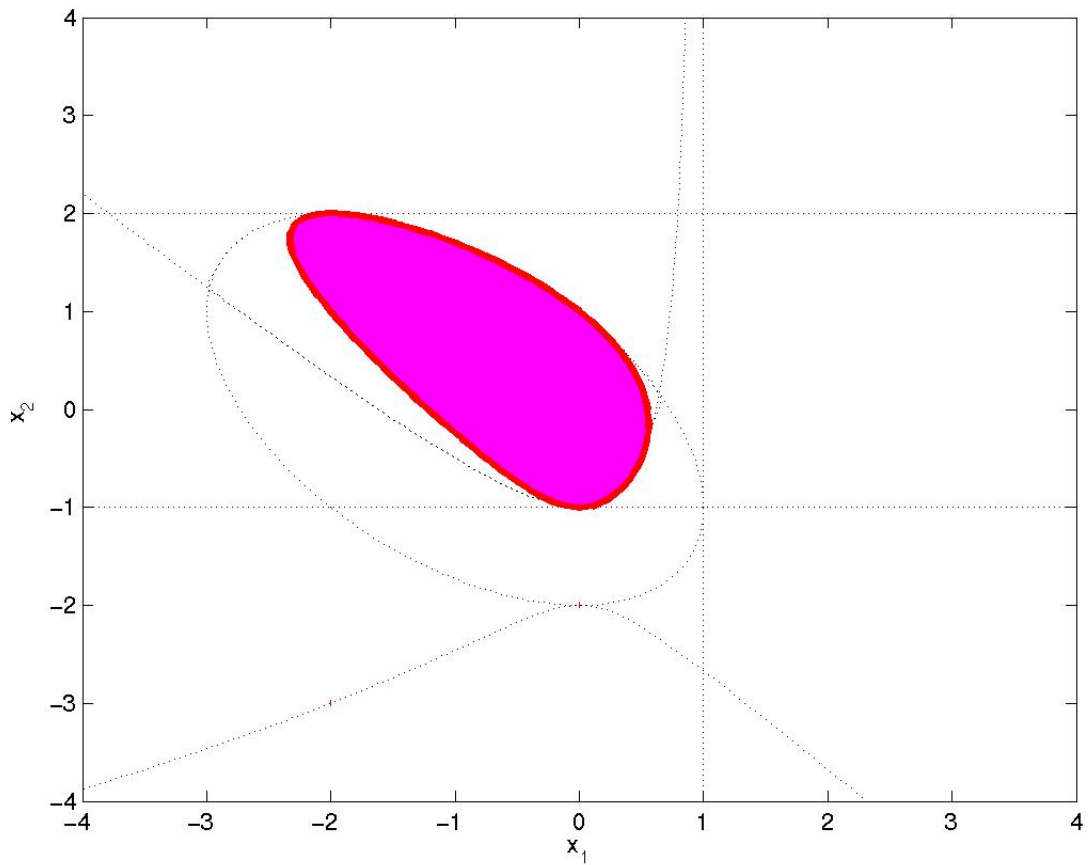


3rd order minor

$$f_7(x) = (1 + x_2)((1 - x_1)(2 - x_2) - (x_1 + x_2)^2) - x_1^2(2 - x_2) \geq 0$$



LMI feasible set = intersection of
semialgebraic sets $f_i(x) \geq 0$
for $i = 1, \dots, 7$

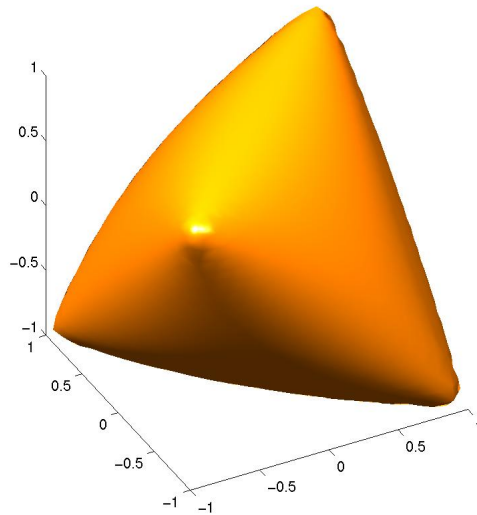


Example of 3D LMI feasible set

LMI set

$$\mathcal{F} = \{x \in \mathbb{R}^3 : \begin{bmatrix} 1 & x_1 & x_2 \\ x_1 & 1 & x_3 \\ x_2 & x_3 & 1 \end{bmatrix} \succeq 0\}$$

arising in SDP relaxation of MAXCUT



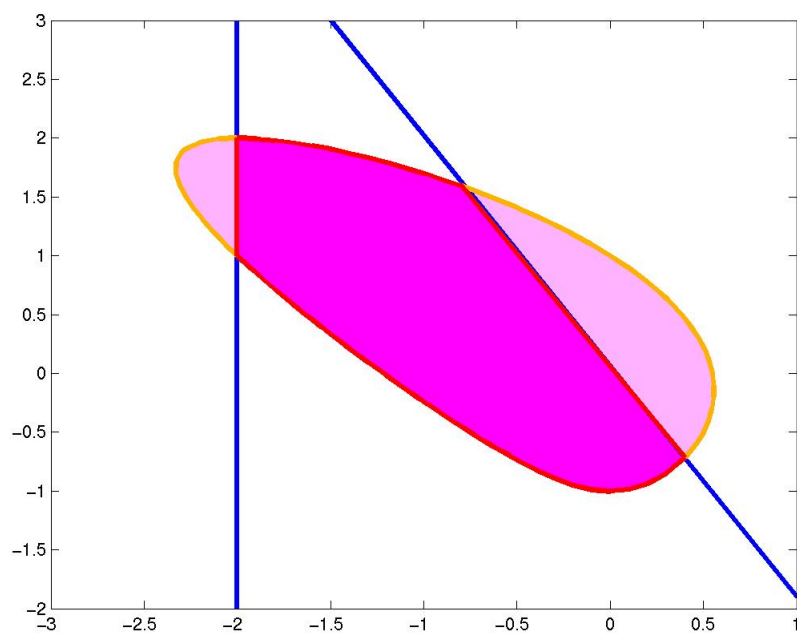
Semialgebraic set

$$\mathcal{F} = \{x \in \mathbb{R}^3 : 1 + 2x_1x_2x_3 - (x_1^2 + x_2^2 + x_3^2) \geq 0, \\ x_1^2 \leq 1, x_2^2 \leq 1, x_3^2 \leq 1\}$$

Intersection of LMI sets

Intersection of LMI feasible sets

$$F(x) \succeq 0 \quad x_1 \geq -2 \quad 2x_1 + x_2 \leq 0$$



is also an LMI

$$\begin{bmatrix} F(x) & 0 & 0 \\ 0 & x_1 + 2 & 0 \\ 0 & 0 & -2x_1 - x_2 \end{bmatrix} \succeq 0$$

SDP representability

LMI sets are **convex semialgebraic** sets.. but are all convex semialgebraic sets **representable** by LMIs ?

We say that a convex set $X \subset \mathbb{R}^n$ is **SDP representable** if there exists an affine mapping $F(x, u)$ such that

$$x \in X \iff \exists u : F(x, u) \succeq 0$$

In words, if X is the **projection** of the solution set of the LMI $F(x, u) \succeq 0$ onto the x -space and u are **additional**, or **lifting** variables

We say that a convex set $X \subset \mathbb{R}^n$ is **LMI representable** if there exists an affine mapping $F(x)$ such that

$$x \in X \iff F(x) \succeq 0$$

In other words, additional variables u are **not allowed**

Similarly, a convex function $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is SDP or LMI representable if its epigraph

$$\mathcal{F} = \{x, t : f(x) \leq t\}$$

is an SDP or LMI representable set

Conic quadratic forms

The Lorentz, or ice-cream cone

$$\mathcal{L} = \left\{ \begin{bmatrix} x \\ t \end{bmatrix} \in \mathbb{R}^{n+1} : \|x\|_2 \leq t \right\}$$

is SDP (LMI) representable as

$$\mathcal{L} = \left\{ \begin{bmatrix} x \\ t \end{bmatrix} : \begin{bmatrix} tI_n & x \\ x^T & t \end{bmatrix} \succeq 0 \right\}$$

As a result, all (convex quadratic) **conic representable** sets are also SDP representable

The SOCP cone is included in the SDP cone

In the sequel we first give a list of conic representable sets (following Ben-Tal and Nemirovski)

Quadratic forms

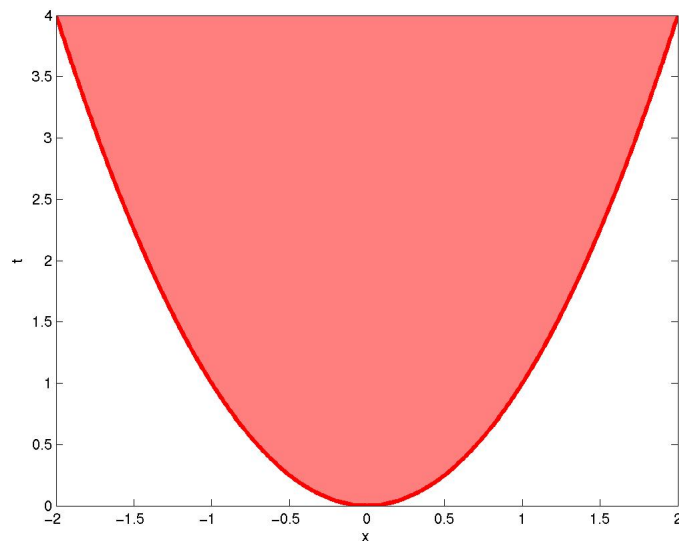
The **Euclidean norm** $\{x, t \in \mathbb{R}^2 : \|x\|_2 \leq t\}$
is conic representable by definition

The **squared Euclidean norm**

$$\{x, t \in \mathbb{R}^2 : x^T x \leq t\}$$

is conic representable as

$$\left\| \begin{bmatrix} x \\ \frac{t-1}{2} \end{bmatrix} \right\|_2 \leq \frac{t+1}{2}$$



Quadratic forms (2)

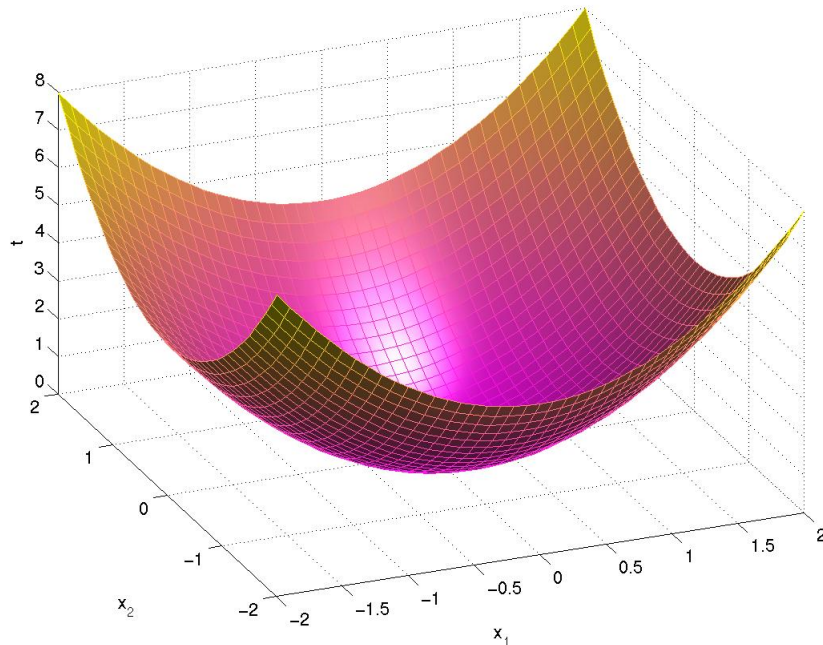
More generally, the **convex quadratic** set

$$\{x \in \mathbb{R}^n, t \in \mathbb{R} : x^T A x + b^T x + c \leq 0\}$$

with $A = A^T \succeq 0$ is conic representable as

$$\left\| \begin{bmatrix} D x \\ \frac{t + b^T x + c}{2} \end{bmatrix} \right\|_2 \leq \frac{t - b^T x - c}{2}$$

where D is the **Cholesky factor** of $A = D^T D$



Who is Cholesky ?

André Louis Cholesky (1875-1918) was a **French military officer** (graduated from Ecole Polytechnique) involved in **geodesy**

He proposed a new procedure for solving **least-squares** triangulation problems

He fell for his country during World War I



Work posthumously published in
Commandant Benoît. Procédé du Commandant Cholesky.
Bulletin Géodésique, No. 2, pp. 67-77, Toulouse,
Privat, 1924.

Nice biography in
C. Brezinski. André Louis Cholesky. *Bulletin of the
Belgian Mathematical Society*, Vol. 3, pp. 45-50, 1996.

I. — NOTICES SCIENTIFIQUES

Commandant BENOIT¹.

NOTE SUR UNE MÉTHODE DE RÉOLUTION DES ÉQUATIONS NORMALES PROVENANT DE L'APPLICATION DE LA MÉTHODE DES MOINDRES CARRÉS A UN SYSTÈME D'ÉQUATIONS LINÉAIRES EN NOMBRE INFÉRIEUR A CELUI DES INCONNUES. — APPLICATION DE LA MÉTHODE A LA RÉOLUTION D'UN SYSTÈME DÉFINI D'ÉQUATIONS LINÉAIRES.

(Procédé du Commandant CHOLESKY².)

Le Commandant d'Artillerie Cholesky, du Service géographique de l'Armée, tué pendant la grande guerre, a imaginé, au cours de recherches sur la compensation des réseaux géodésiques, un procédé très ingénieux de résolution des équations dites *normales*, obtenues par application de la méthode des moindres carrés à des équations linéaires en nombre inférieur à celui des inconnues. Il en a conclu une méthode générale de résolution des équations linéaires.

Nous suivrons, pour la démonstration de cette méthode, la progression même qui a servi au Commandant Cholesky pour l'imaginer.

1. De l'Artillerie coloniale, ancien officier géodésien au Service géographique de l'Armée et au Service géographique de l'Indo-Chine, Membre du Comité national français de Géodésie et Géophysique.

2. Sur le Commandant Cholesky, tué à l'ennemi le 31 août 1918, voir la notice biographique insérée dans le volume du *Bulletin géodésique* de 1922 intitulé : *Union géodésique et géophysique internationale, Première Assemblée générale, Rome, mai 1922, Section de Géodésie*, Toulouse, Privat, 1922, in-8°, 241 p., pp. 159 à 161.

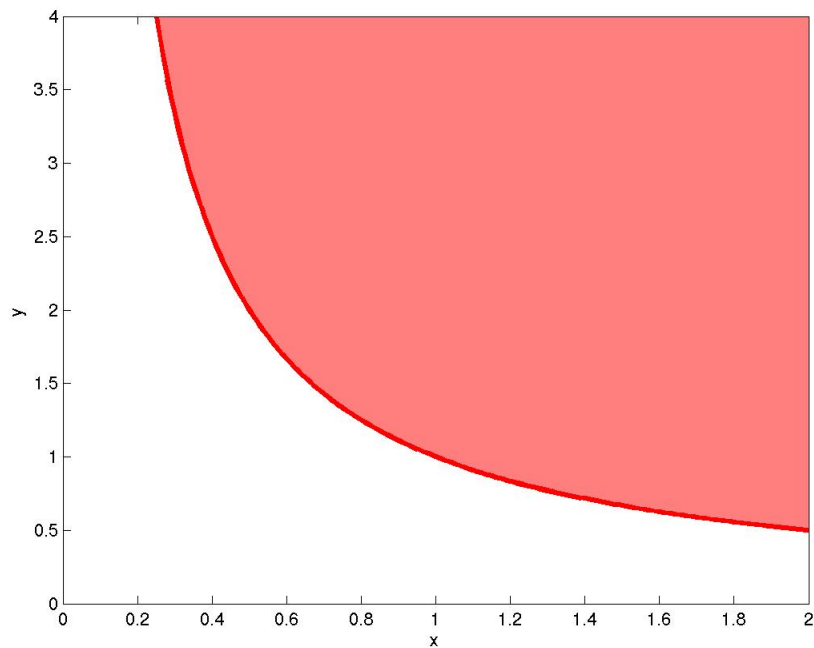
Hyperbola

The branch of **hyperbola**

$$\{x, y \in \mathbb{R}^2 : xy \geq 1, x > 0\}$$

is conic representable as

$$\left\| \begin{bmatrix} \frac{x-y}{2} \\ 1 \end{bmatrix} \right\|_2 \leq \frac{x+y}{2}$$



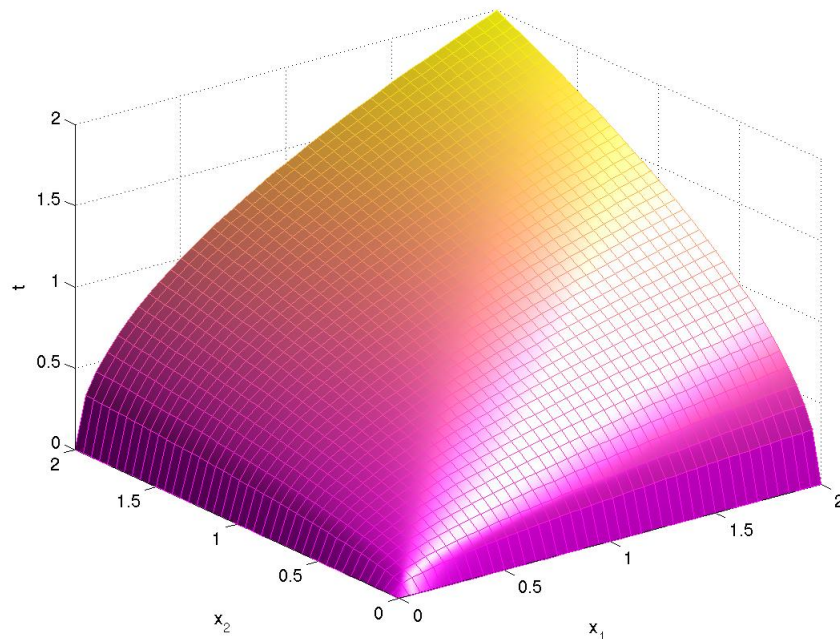
Geometric mean of two variables

The hypograph of the **geometric mean** of 2 variables

$$\{x_1, x_2, t \in \mathbb{R}^3 : x_1, x_2 \geq 0, \sqrt{x_1 x_2} \geq t\}$$

is conic representable as

$$\exists u : u \geq t, \left\| \begin{bmatrix} u \\ \frac{x_1 - x_2}{2} \end{bmatrix} \right\|_2 \leq \frac{x_1 + x_2}{2}$$



Geometric mean of several variables

The hypograph of the geometric mean of 2^k variables

$$\{x_1, \dots, x_{2^k}, t \in \mathbb{R}^{2^k+1} : x_i \geq 0, \sqrt{x_1 \cdots x_{2^k}} \geq t\}$$

is also conic representable

Proof: [iterate](#) the previous construction

Example with $k = 3$:

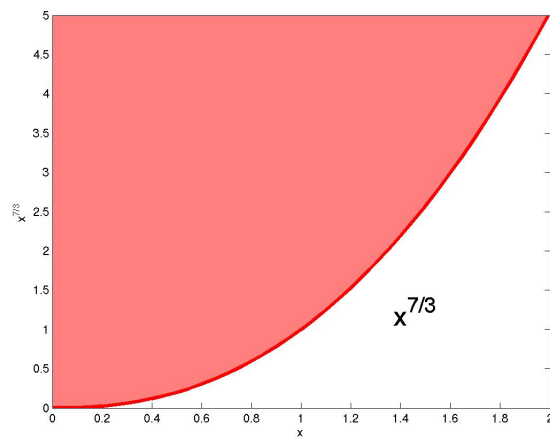
$$\begin{array}{l} (x_1 x_2 \cdots x_8)^{1/8} \geq t \\ \left. \begin{array}{l} \sqrt{x_1 x_2} \geq x_{11} \\ \sqrt{x_3 x_4} \geq x_{12} \\ \sqrt{x_5 x_6} \geq x_{13} \\ \sqrt{x_7 x_8} \geq x_{14} \end{array} \right\} \left. \begin{array}{l} \sqrt{x_{11} x_{12}} \geq x_{21} \\ \sqrt{x_{13} x_{14}} \geq x_{22} \end{array} \right\} \sqrt{x_{21} x_{22}} \geq x_{31} \geq t \end{array}$$

Useful idea in other SDP representability problems

Rational functions

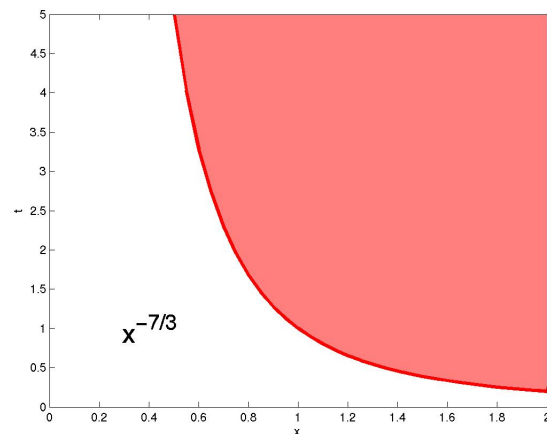
Usually similar ideas, we can show that the increasing **rational power functions**

$$f(x) = x^{p/q}, \quad x \geq 0$$



with rational $p/q \geq 1$, as well as the decreasing

$$g(x) = x^{-p/q}, \quad x \geq 0$$



with rational $p/q \geq 0$, are both conic representable

Rational functions

Example:

$$\{x, t : x \geq 0, x^{7/3} \leq t\}$$

Start from conic representable

$$\hat{t} \leq (\hat{x}_1 \cdots \hat{x}_8)^{1/8}$$

and replace

$$\begin{aligned}\hat{t} &= \hat{x}_1 = x \geq 0 \\ \hat{x}_2 &= \hat{x}_3 = \hat{x}_4 = t \geq 0 \\ \hat{x}_5 &= \hat{x}_6 = \hat{x}_7 = \hat{x}_8 = 1\end{aligned}$$

to get

$$\begin{aligned}x &\leq x^{1/8}t^{3/8} \\ x^{7/8} &\leq t^{3/8} \\ x^{7/3} &\leq t\end{aligned}$$

Same idea works for any rational $p/q \geq 1$

- **lift** = use additional variables, and
- **project** in the space of original variables

Even power monomial

The epigraph of **even power monomial**

$$\mathcal{F} = \{x, t : x^{2p} \leq t\}$$

where p is a positive integer
is conic representable

Note that

$$\{x, t : x^{2p} \leq t\}$$

$$\iff$$

$$\begin{aligned} & \{x, y, t : x^2 \leq y\} \\ & \{x, y, t : y \geq 0, y^p \leq t\} \end{aligned}$$

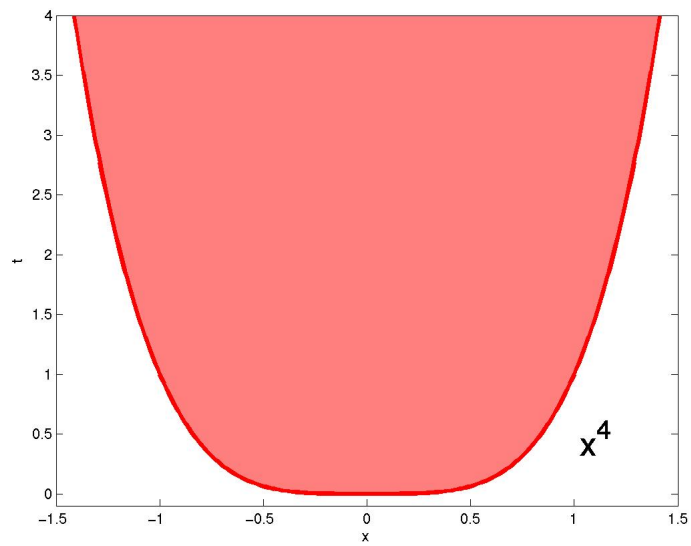
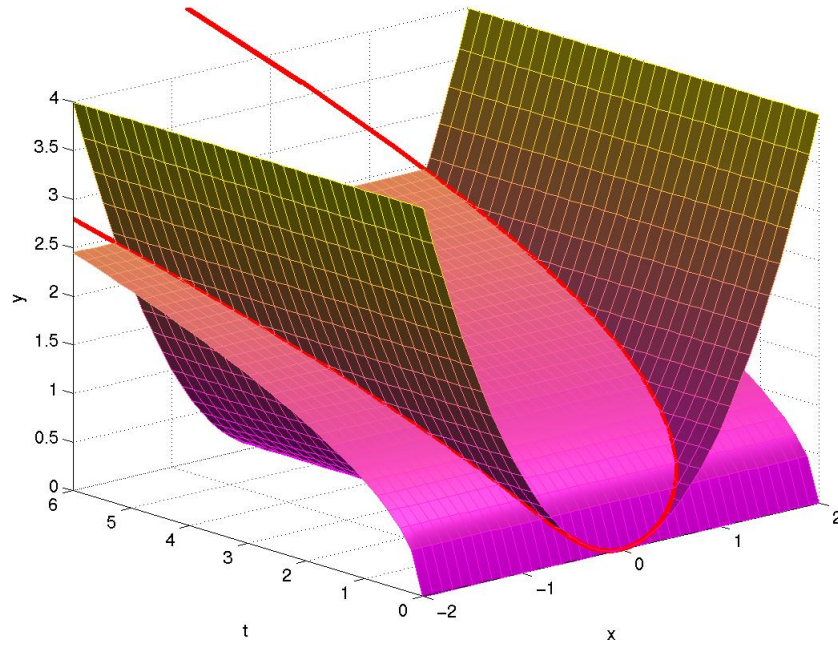
both conic representable

Use **lifting** y and **project** back onto x, t

Similarly, **even power polynomials** are conic representable (combinations of monomials)

Even power monomial

$$\mathcal{F} = \{x, t : x^4 \leq t\}$$



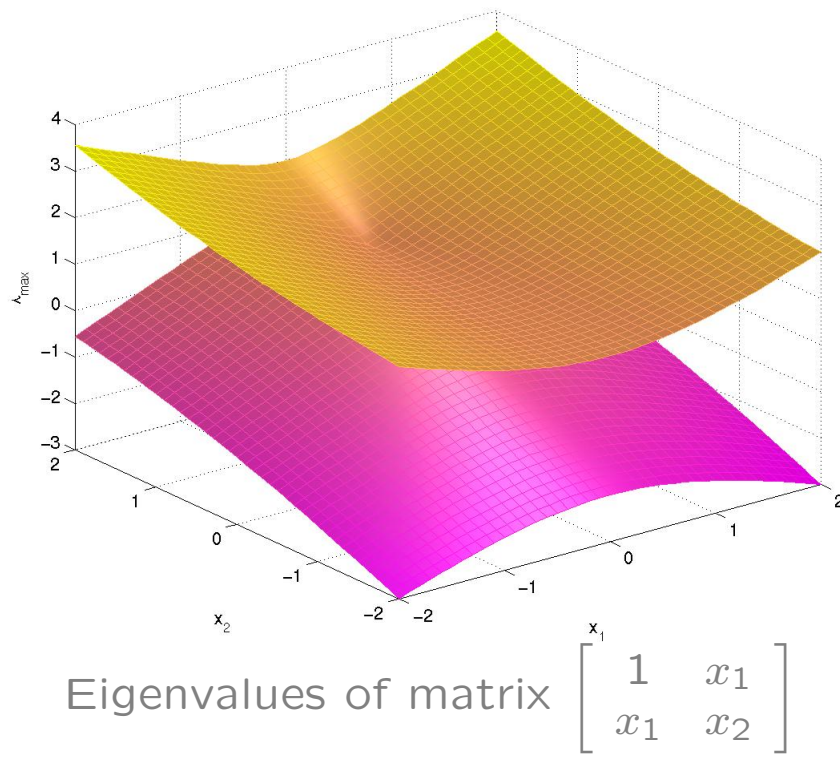
Largest eigenvalue

The epigraph of the function **largest eigenvalue** of a symmetric matrix

$$\{X = X^T \in \mathbb{R}^{n \times n}, t \in \mathbb{R} : \lambda_{\max}(X) \leq t\}$$

is SDP (LMI) representable as

$$X \preceq tI_n$$



Sums of largest eigenvalues

Let

$$S_k(X) = \sum_{i=1}^k \lambda_i(X), \quad k = 1, \dots, n$$

denote the **sum of the k largest eigenvalues** of an n -by- n symmetric matrix X

The epigraph

$$\{X = X^T \in \mathbb{R}^{n \times n}, t \in \mathbb{R} : S_k(X) \leq t\}$$

is SDP representable as

$$\begin{aligned} t - ks - \text{trace } Z &\succeq 0 \\ Z &\succeq 0 \\ Z - X + sI_n &\succeq 0 \end{aligned}$$

where Z and s are additional variables

Determinant of a PSD matrix

The **determinant**

$$\det(X) = \prod_{i=1}^n \lambda_i(X)$$

is not a convex function of X , but the function

$$f_q(X) = -\det^q(X), \quad X = X^T \succeq 0$$

is convex when $q \in [0, 1/n]$ is rational

The epigraph

$$\{f_q(X) \leq t\}$$

is SDP representable as

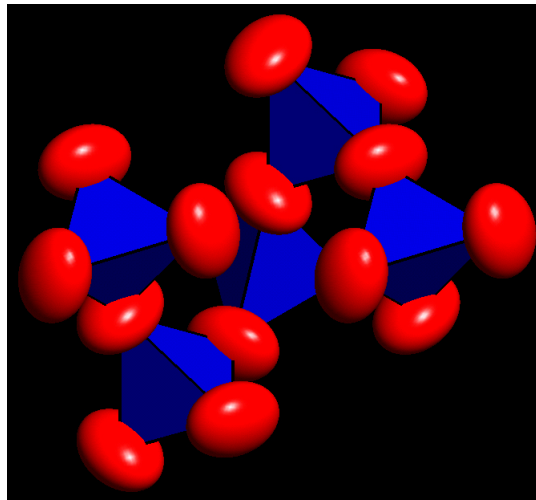
$$\begin{bmatrix} X & \Delta \\ \Delta^T & \text{diag } \Delta \\ t \leq (\delta_1 \cdots \delta_n)^q \end{bmatrix} \succeq 0$$

since we know that the latter constraint (hypograph of a concave monomial) is conic representable

Here Δ is a lower triangular matrix of additional variables with diagonal entries δ_i

Application: extremal ellipsoids

A little excursion in the world of ellipsoids and polytopes..



Crystal structure

Various representations of an ellipsoid in \mathbb{R}^n

$$\begin{aligned} E &= \{x \in \mathbb{R}^n : x^T P x + 2x^T q + r \leq 0\} \\ &= \{x \in \mathbb{R}^n : (x - x_c)^T P (x - x_c) \leq 1\} \\ &= \{x = Qy + x_c \in \mathbb{R}^n : y^T y \leq 1\} \\ &= \{x \in \mathbb{R}^n : \|Rx - x_c\| \leq 1\} \end{aligned}$$

where

$$Q = R^{-1} = P^{-1/2} \succ 0$$

Ellipsoid volume

Volume of ellipsoid $E = \{Qy + x_c : y^T y \leq 1\}$

$$\text{vol } E = k_n \det Q$$

where k_n is volume of n -dimensional unit ball

$$k_n = \begin{cases} \frac{2^{(n+1)/2} \pi^{(n-1)/2}}{n(n-2)!!} & \text{for } n \text{ odd} \\ \frac{2\pi^{n/2}}{n(n/2-1)!} & \text{for } n \text{ even} \end{cases}$$

n	1	2	3	4	5	6	7	8
k_n	2.00	3.14	4.19	4.93	5.26	5.17	4.72	4.06

Unit ball has maximum volume for $n = 5$!

Outer and inner ellipsoidal approximations

Let $S \subset \mathbb{R}^n$ be a **solid** = a closed bounded convex set with nonempty interior

- the largest volume ellipsoid E_{in} contained in S is unique and satisfies

$$E_{\text{in}} \subset S \subset nE_{\text{in}}$$

- the smallest volume ellipsoid E_{out} containing S is unique and satisfies

$$E_{\text{out}}/n \subset S \subset E_{\text{out}}$$

These are **Löwner-John** ellipsoids

Factor n reduces to \sqrt{n} if S is symmetric

How can these ellipsoids be computed ?

Ellipsoid in polytope

Let the intersection of hyperplanes

$$S = \{x \in \mathbb{R}^n : a_i^T x \leq b_i, i = 1, \dots, m\}$$

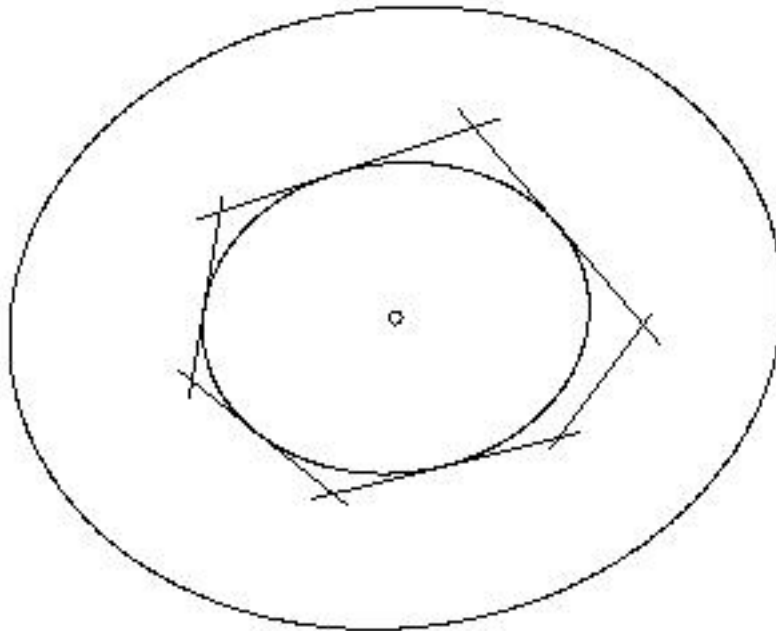
describe a **polytope** = bounded nonempty polyhedron

The **largest volume ellipsoid contained** in S is

$$E = \{Qy + x_c : y^T y \leq 1\}$$

where Q, x_c are optimal solutions of the LMI

$$\begin{aligned} \max \quad & \det^{1/n} Q \\ & Q \succeq 0 \\ & \|Qa_i\|_2 \leq b_i - a_i^T x_c, \quad i = 1, \dots, m \end{aligned}$$



Polytope in ellipsoid

Let the convex hull of vertices

$$S = \text{conv} \{x_1, \dots, x_m\}$$

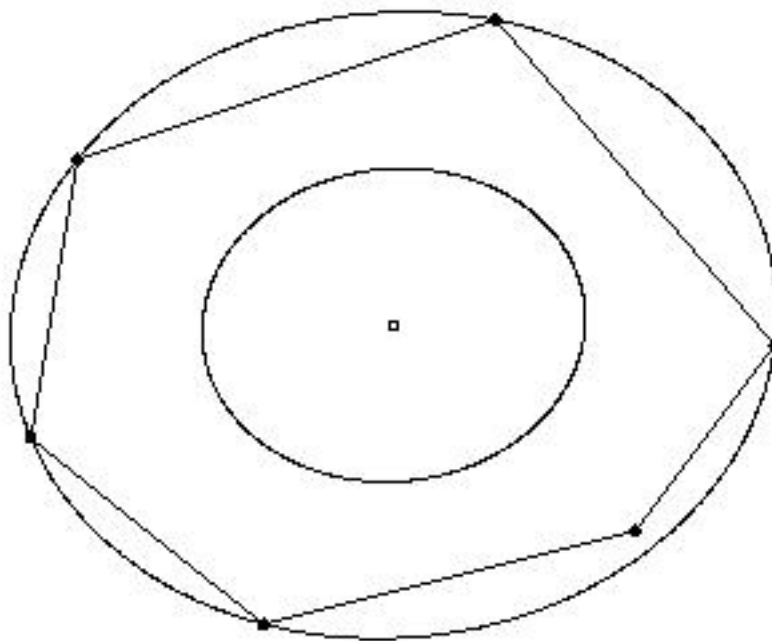
describe a polytope

The smallest volume ellipsoid containing S is

$$E = \{x : (x - x_c)^T P (x - x_c) \leq 1\}$$

where $P, x_c = -P^{-1}q$ are optimal solutions of the LMI

$$\begin{aligned} \max \quad & t \\ & t \leq \det^{1/n} P \\ & \begin{bmatrix} P & q \\ q^T & r \end{bmatrix} \succeq 0 \\ & x_i^T P x_i + 2x_i^T q + r \leq 1, \quad i = 1, \dots, m \end{aligned}$$



Sums of largest singular values

Let

$$\Sigma_k(X) = \sum_{i=1}^k \sigma_i(X), \quad k = 1, \dots, n$$

denote the **sum of the k largest singular values** of an n -by- m matrix X

Then the epigraph

$$\{X = X^T \in \mathbb{R}^{n \times n}, t \in \mathbb{R} : \Sigma_k(X) \leq t\}$$

is SDP representable since

$$\sigma_i(X) = \lambda_i \left(\begin{bmatrix} 0 & X^T \\ X & 0 \end{bmatrix} \right)$$

and the sum of largest eigenvalues of a symmetric matrix is SDP representable

Schur complement

We can use the **Schur complement** to convert a non-linear matrix inequality into an LMI

$$\begin{bmatrix} A(\mathbf{x}) & B(\mathbf{x}) \\ B^T(\mathbf{x}) & C(\mathbf{x}) \end{bmatrix} \succeq 0 \iff \begin{matrix} A(\mathbf{x}) - B(\mathbf{x})C^{-1}(\mathbf{x})B^T(\mathbf{x}) \succeq 0 \\ C(\mathbf{x}) \succ 0 \end{matrix}$$



Issai Schur
(1875 Mogilyov - 1941 Tel Aviv)

Elimination lemma

To remove decision variables we can use the [elimination lemma](#)

$$\begin{aligned} A(\boldsymbol{x}) + B(\boldsymbol{x})X C^T(\boldsymbol{x}) + C(\boldsymbol{x})X^T B^T(\boldsymbol{x}) > 0 \\ \iff \\ \tilde{B}^T(\boldsymbol{x})A(\boldsymbol{x})\tilde{B}(\boldsymbol{x}) > 0 \quad \tilde{C}^T(\boldsymbol{x})A(\boldsymbol{x})\tilde{C}(\boldsymbol{x}) > 0 \end{aligned}$$

where \tilde{B} and \tilde{C} are orthogonal complements of B and C respectively, and \boldsymbol{x} is a decision variable independent of matrix X

Can be shown with SDP duality..

Particular case: [Finsler's theorem](#)

Positive polynomials

The set of univariate polynomials that are positive on the real axis is a **convex** set that can be described by an LMI

Can be proved with cone duality (Nesterov) or with theory of moments (Lasserre)

The even polynomial

$$p(s) = p_0 + p_1s + \cdots + p_{2n}s^{2n}$$

satisfies $p(s) \geq 0$ for all $s \in \mathbb{R}$ if and only if

$$\begin{aligned} p_k &= \sum_{i+j=k} X_{ij}, & k &= 0, 1, \dots, 2n \\ &= \text{trace } H_k X \end{aligned}$$

for some matrix $X = X^T \succeq 0$

Sum-of-squares decomposition

The expression of p_k with Hankel matrices H_k comes from

$$p(s) = [1 \quad s \quad \dots \quad s^n] X [1 \quad s \quad \dots \quad s^n]^*$$

hence $X \succeq 0$ naturally implies $p(s) \geq 0$

Conversely, existence of X for any polynomial $p(s) \geq 0$ follows from the existence of a **sum-of-squares** decomposition (with at most two elements) of

$$p(s) = \sum_k q_k^2(s) \geq 0$$

Matrix X has entries $X_{ij} = \sum_k q_{k_i} q_{k_j}$

Primal and dual formulations

Global minimization of polynomial

$$p(s) = \sum_{k=0}^n p_k s^k$$

Global optimum p^* : maximum value of \hat{p} such that $p(s) - \hat{p}$ stays globally nonnegative

Primal LMI

$$\begin{aligned} \max \quad & \hat{p} = p_0 - \text{trace } H_0 \mathbf{X} \\ \text{s.t.} \quad & \text{trace } H_k \mathbf{X} = p_k, \quad k = 1, \dots, n \\ & \mathbf{X} \succeq 0 \end{aligned}$$

Dual LMI

$$\begin{aligned} \min \quad & p_0 + \sum_{k=1}^n p_k y_k \\ \text{s.t.} \quad & H_0 + \sum_{k=1}^n H_k y_k \succeq 0 \end{aligned}$$

with **Hankel** structure (moment matrix)

Positive polynomials and LMIs

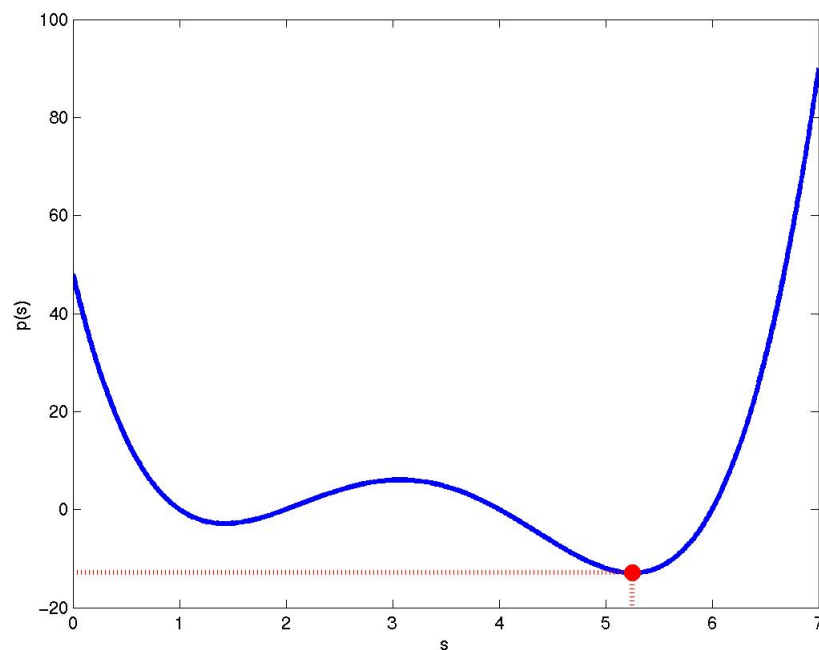
Example: **Global minimization** of the polynomial

$$p(s) = 48 - 92s + 56s^2 - 13s^3 + s^4$$

We just have to solve the **dual LMI**

$$\begin{array}{ll} \min & 48 - 92y_1 + 56y_2 - 13y_3 + y_4 \\ \text{s.t.} & \begin{bmatrix} 1 & y_1 & y_2 \\ y_1 & y_2 & y_3 \\ y_2 & y_3 & y_4 \end{bmatrix} \succeq 0 \end{array}$$

to obtain $p^* = p(5.25) = -12.89$



Complex LMIs

The **complex** valued LMI

$$F(\mathbf{x}) = A(\mathbf{x}) + jB(\mathbf{x}) \succeq 0$$

is equivalent to the real valued LMI

$$\begin{bmatrix} A(\mathbf{x}) & B(\mathbf{x}) \\ -B(\mathbf{x}) & A(\mathbf{x}) \end{bmatrix} \succeq 0$$

If there is a complex solution to the LMI
then there is a **real** solution to the same LMI

Note that matrix $A(\mathbf{x}) = A^T(\mathbf{x})$ is symmetric
whereas $B(\mathbf{x}) = -B^T(\mathbf{x})$ is skew-symmetric

Rigid convexity

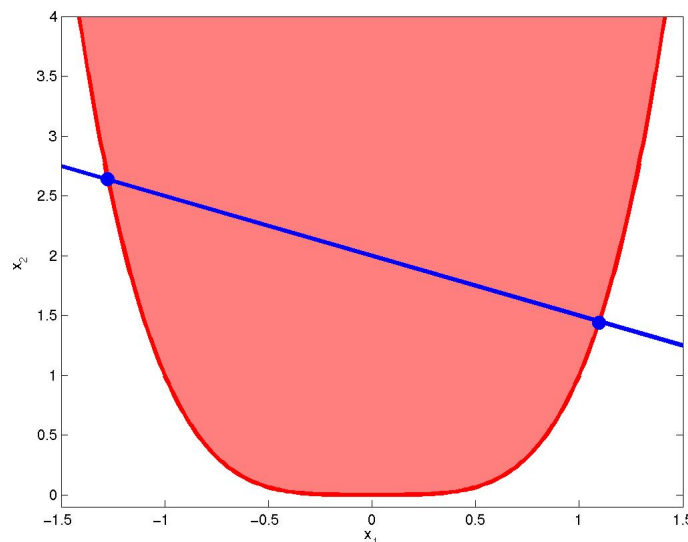
Helton & Vinnikov showed that a convex 2D set

$$\mathcal{F} = \{x \in \mathbb{R}^2 : p(x) \geq 0\}$$

defined by a polynomial $p(x)$ of minimum degree d is LMI representable **without lifting variables** iff \mathcal{F} is **rigidly convex**, meaning that

for every point $x \in X$ and almost every line through x then the line intersects $p(x) = 0$ in exactly d points

Example: $\mathcal{F} = \{x_1, x_2 \in \mathbb{R}^2 : p(x) = x_2 - x_1^4 \geq 0\}$
with 2 line intersections
is not rigidly convex because $2 < d = 4$



.. but it is LMI representable **with lifting variables**
see the previous construction for even power monomials