

The RO-6 Theorem and Latency-Optimal Read-Only Transactions



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Outline

- Read-only transactions
- Latency-optimality
- The RO-6 theorem
- New algorithms
- Preliminary results

Target: Scalable Data Stores

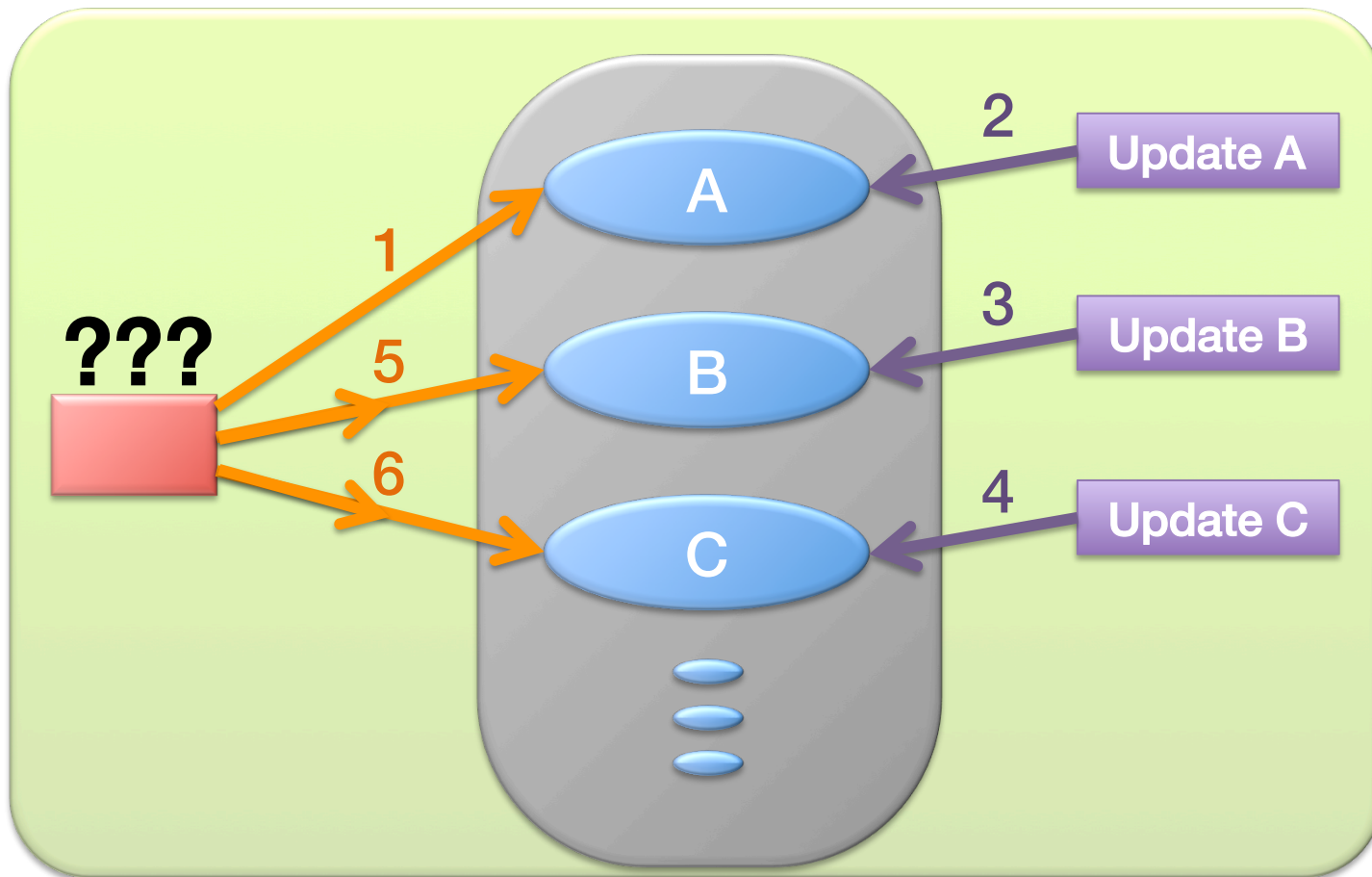


Key Point: Data is distributed across multiple nodes

Must contact multiple disjoint servers to read data

Read-Only Txns Necessary

Asynchronous requests + distributed data = ??????



Why Latency?

- **Directly impacts user experience**
- **Multiplicative**
 - “In practice, a single user request [has] critical path that is dozens of subqueries long.” [Facebook, HotOS ‘15]
- **Matters even more for geo-partitioned data**
 - Khiem Ngo “K2” project looking at this
- **Tractable!**
 - Throughput is much harder to capture precisely

The RO-6 Theorem:

Impossibility Result for All 6 Properties

1. One round of messages
 2. One version returned by servers
 3. Non-blocking
 4. Strict serializability
 5. Conflicting write transactions
 6. Point-to-point messages
- Latency
- Power
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- A diagram consisting of two large right-facing curly braces. The top brace groups items 1, 2, and 3, with the word 'Latency' to its right. The bottom brace groups items 4 and 5, with the word 'Power' to its right. Item 6 is not grouped by either brace.

RO-6 Theorem

- Impossible to achieve all 6 together
 - Proof by contradiction
- Set of 6 properties is minimal
 - Any 5 of the 6 properties is achievable
 - We have the algorithms!
 - Still determining exact bounds on properties
 - 1 rounds vs 2 rounds vs N rounds

Why RO-6 Theorem?

- **Avoid trying to do the impossible**
 - This is how we got here
- **Let us actually say this is the best we can do**
 - Caveat: latency + read-only transactions
- **Guide for understanding when we can improve read-only transaction algorithms**

New Algorithms

Make reads faster at the expense of writes

- **COPS-RO6: No conflicting write-only transactions**
 - Writes check causal dependencies before being applied to determine preceding read-only txns
- **Eiger-RO6: Blocking**
 - 2PC for write transactions carries metadata about “earlier” read-only transactions, read-only transactions block to eliminate race-condition
- **Eiger-Chris: Strict serializability**
 - Provides “process serializability”, delays write completion
- **Rococo-RO6: Blocking**

Implementation & Evaluation

- COPS-RO6: Still to be implemented
- Eiger-RO6: Implemented
 - Lower latency, especially for large “expanding” read-only txns
 - No throughput improvement
 - Conjecture: throughput will improve for highly skewed workloads
- Eiger-Chris: Still to be implemented
- Rococo-RO6: Implemented
 - Lower latency & higher throughput
 - Rococo eliminated aborts for read/write txns, but read-only txns could abort. New algorithm for read-only txns avoid all aborts.
 - TPC-C: All 5 types of transactions avoid aborts now

Feedback

- **New targets for RO-6 Theorem?**
- **Related work pointers**
- **Better terminology**
- **Eliminate “point-to-point” messages**