

Towards Quantitative Security Evaluation ?

Yves Deswarte,
Mohamed Kaâniche

LAAS-CNRS
Toulouse, France



Fault forecasting

= Evaluation:

- Gain confidence that system dependability is satisfactory
- Select architecture/components to achieve the best **dependability-performance-cost** trade-off

❖ Quantitative measures

- Reliability: MTFF = mean time to first failure,
 $R(t) = \text{prob}_{\text{continuous service}}(t)$
- Availability: $MTBF / (MTBF + MTTR)$,
 $A(t) = \text{prob}_{\text{correct service provided when needed}}(t)$

Basic assumption

- ❖ Faults = elementary component failures
(or other rare physical phenomena)
Model = independent stochastic processes
with known distributions
- ❖ OK for physical H/W faults
and most environmental faults
- ❖ ~OK for most S/W design faults (bugs)
- ❖ **Not OK** for attacks or malicious design faults

Security Evaluation

❖ Usual techniques

- Evaluation criteria (TCSEC, ITSEC, CC, ...):
~ qualitative evaluation
- Risk assessment: subjective evaluation of vulnerabilities, threats, consequences
- These are static analyses rather than dynamic:
"How the system has been built?" rather than
"How is it operated?"

Quantitative security evaluation

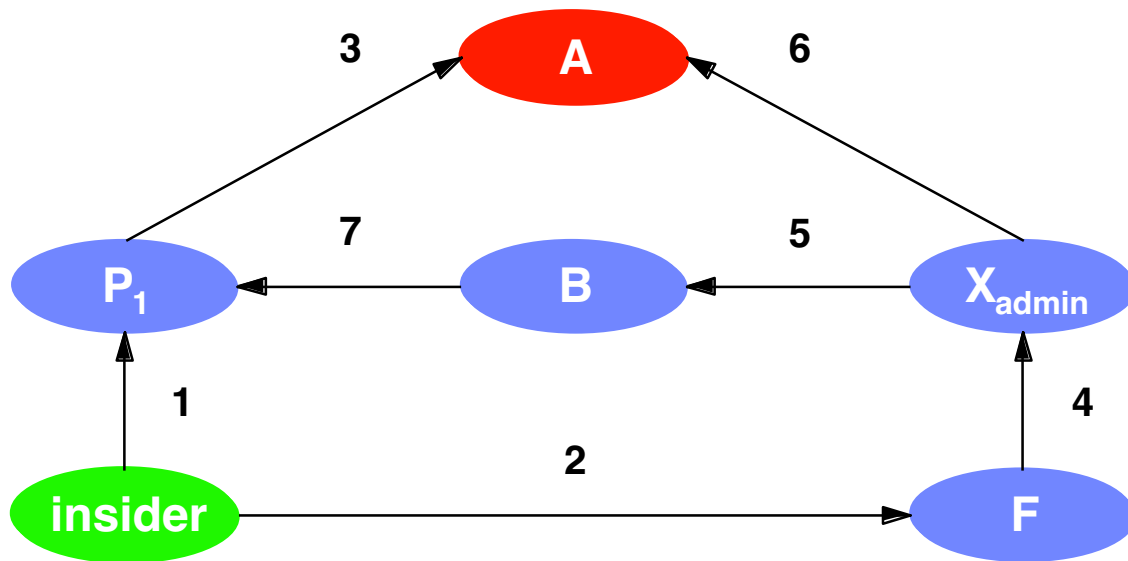
- ❖ Measure = **effort** needed for a possible attacker to defeat the security policy
- ❖ Objectives:
 - Take into account security/usability trade-offs
 - Monitor security evolutions according to configuration and use changes
 - Identify the best security improvement for the least usability change

ESOPE: General approach

- ❖ Identify security objectives: security policy
- ❖ Model (operational) system vulnerabilities
- ❖ Model the attack processes
- ❖ Compute significant measures

Vulnerability modeling

Privilege graph



- 1) X can guess Y's password
- 2) X can install a Trojan horse that Y can activate
- 3) X can exploit a flaw in Y's mailer
- 4) Y is a subset of X
- 5) Y uses a program that X can modify
- 6) X can modify a "s-uid" program owned by Y
- 7) X is in Y's .rhosts

- ❖ **Node** = a set of privileges (user, group, rôle, ...)
- ❖ **Arc** = a method to transfer privileges = vulnerability
- ❖ **Path** = set of vulnerabilities usable by a possible attacker to reach a target
- ❖ **Weight** = for each arc, effort to exploit the arc's vulnerability

Assumptions on the attack process

- ❖ Attack process = all possible successful attack scenarios
- ❖ Reasonable assumptions
 - The attacker knows only the vulnerabilities that can be exploited with the privileges he already owns.
 - The attacker will not exploit vulnerabilities which would give him privileges he already owns.
- ❖ Plus one out of the two following assumptions:
 - *Total Memory (TM)*: the attacker remembers all the vulnerabilities he did not exploit in the previous steps, and he can "back-track".
 - *Memory-Less (ML)*: the attacker considers only the vulnerabilities that can be exploited with the new privileges he just acquired.

Measure computation

① Identify the attacker-target couples

② For each couple, compute:

METF-ML: Mean Effort To security Failure
(i.e. to reach the target) with ML assumption.

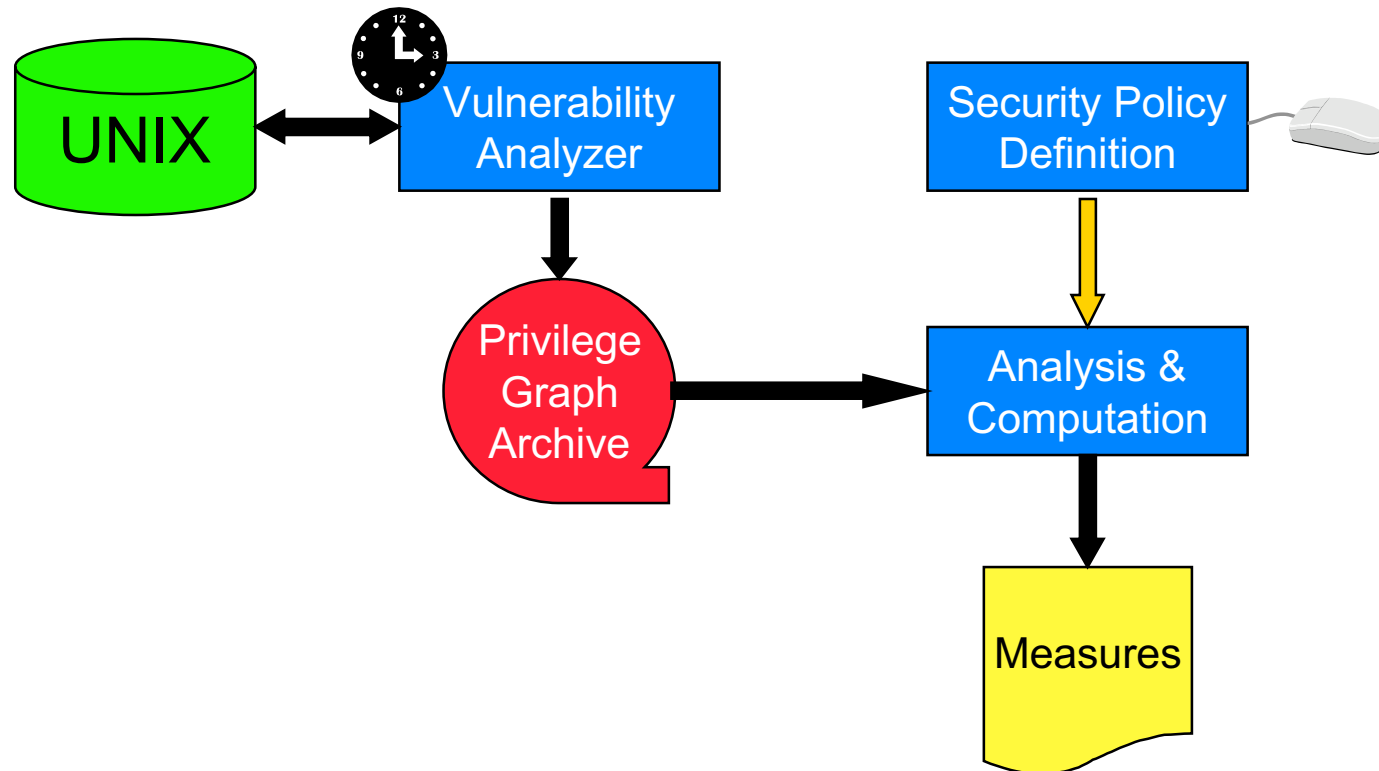
METF-TM: Mean Effort To security Failure with TM
assumption.

Shortest Path: Mean effort to go through the shortest path.

Number of Paths: Number of possible paths from the attacker
to the target nodes.

ESOPE Tool Set

(Evaluation de la Sécurité OPÉrationnelle)



Experiment report

❖ Objectives:

○ Validate the approach:

- Assess the measure pertinence wrt. system changes (configuration, users, ...)
- Feasibility of a full-size system evaluation.

○ *Was not aimed:*

- Correct the identified vulnerabilities

Experiment context

Target system:

- Unix
- 700 users -
300 machines - LAN
- 13 months
(June 1995 - July 1996)

13 types of vulnerabilities
(files `.rhosts`, `.*rc`, passwords, etc.)

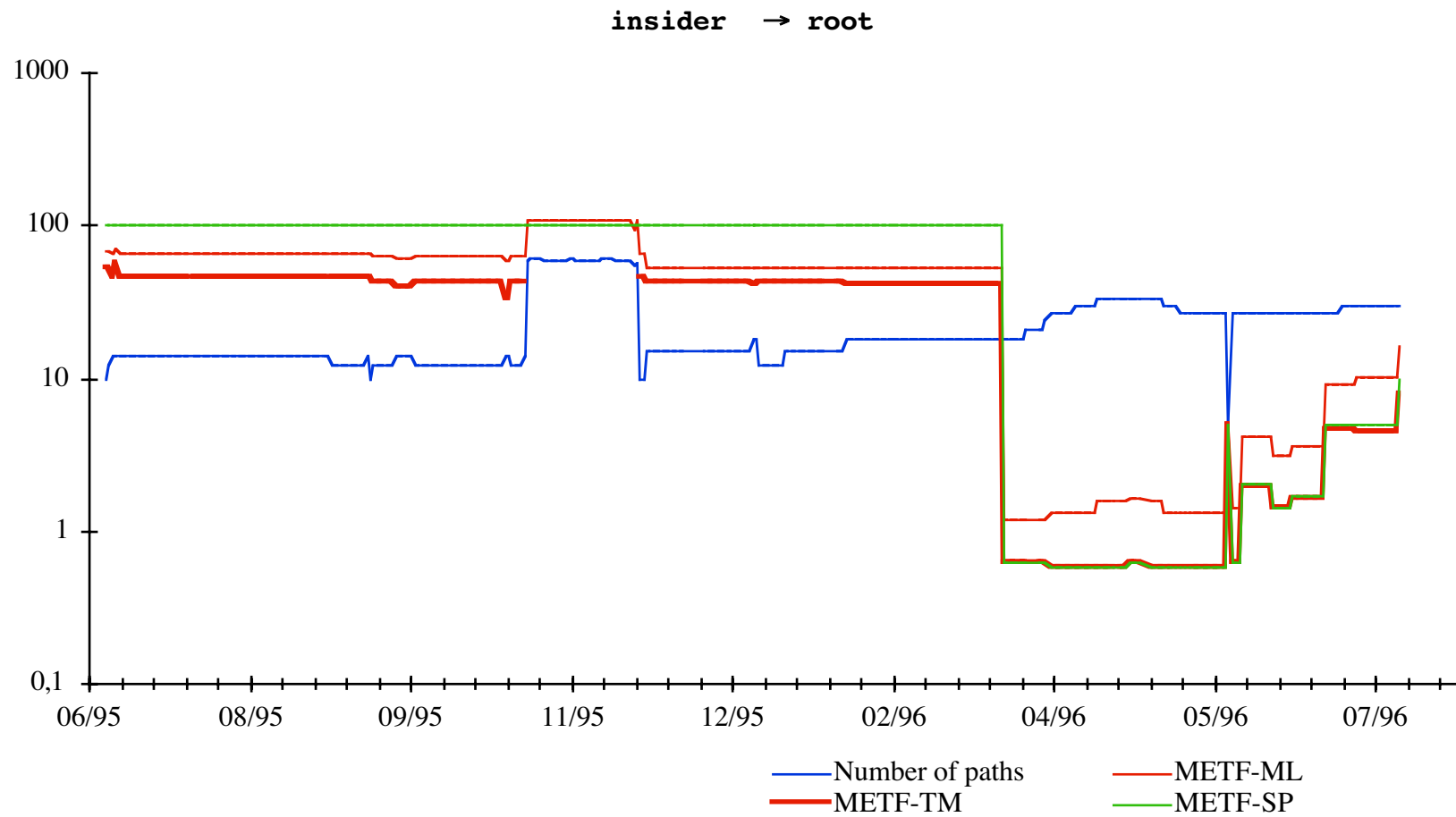
Security objectives:

	Attacker	Target
Objective 1	insider	root
Objective 2	insider	admin_group

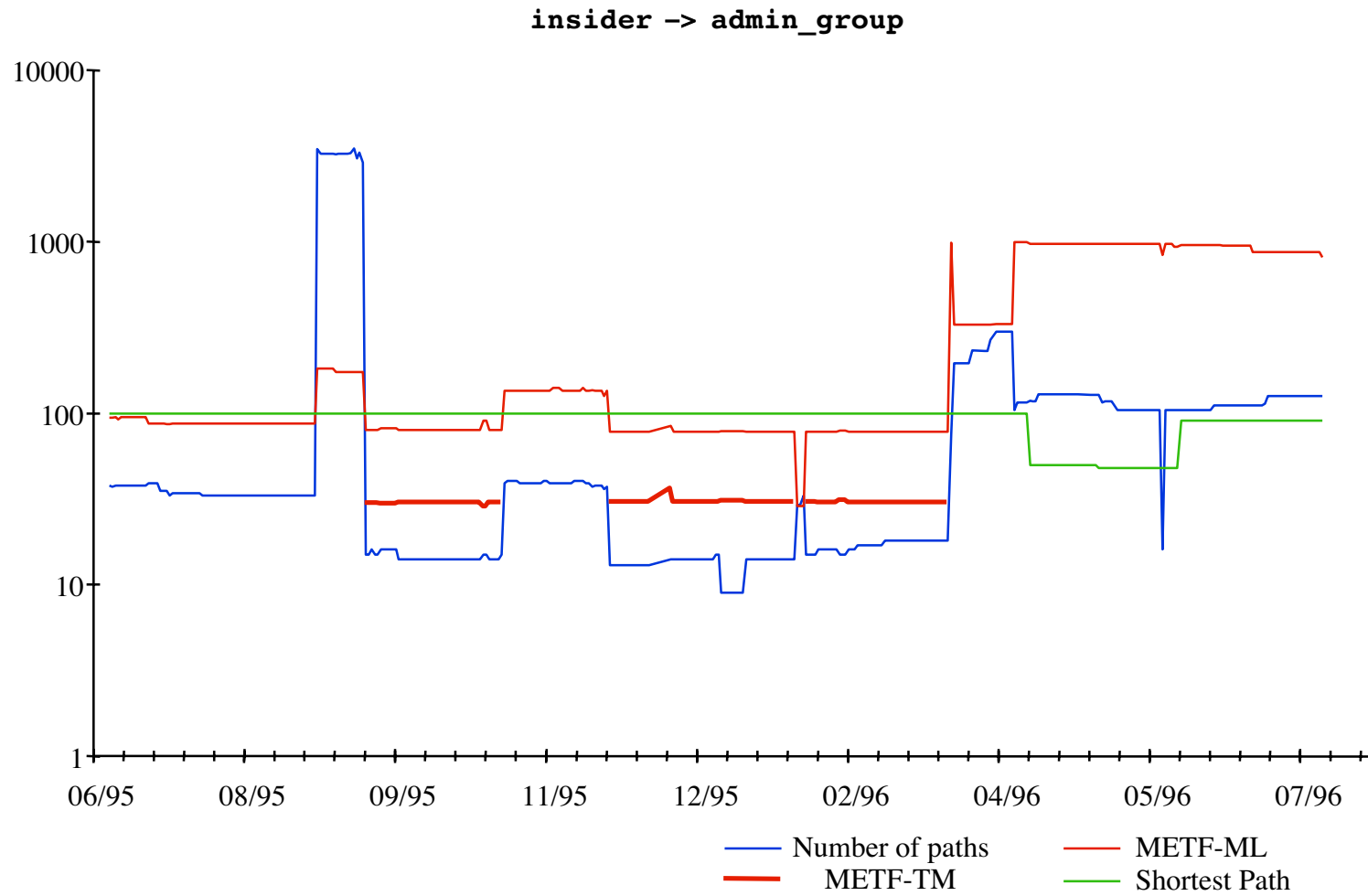
4 difficulty levels:

Type	Weight
immediate	10
easy	10^2
difficult	10^3
very difficult	10^4

Results (1)



Results (2)



Comparison between measures

- ❖ The shortest path (**SP**) is not sensitive enough to identify important events
- ❖ The number of paths (**NP**) changes too often and would produce a large number of false alarms.
- ❖ **METF-ML** presents a good sensitivity to important events.
- ❖ **METF-TM** is easier to interpret, but is sometimes too complex to be computed.

Problems

- ❖ Is the model valid in the real world?
 - ❖ TM and ML are 2 extreme attacker behaviors, but what would be a "real" attacker behavior?
 - ❖ Weight parameters are assessed arbitrarily (**subjective** ?)
 - ❖ Tenacity? Collusion? Attack rates?
- We need real data!

Validation based on real attack data

- ❖ Collect real life data to learn and analyze attackers behaviors, tools and tactics
- ❖ Objectives
 - Validate attack assumptions
 - Analyze adequacy of the privilege graph to describe new vulnerabilities and derive attack scenarios
 - Extend security evaluation approach by taking into account distribution of attacks in time, correlation between attacks, etc.

Honeypots and Honeynets

❖ Honeypot

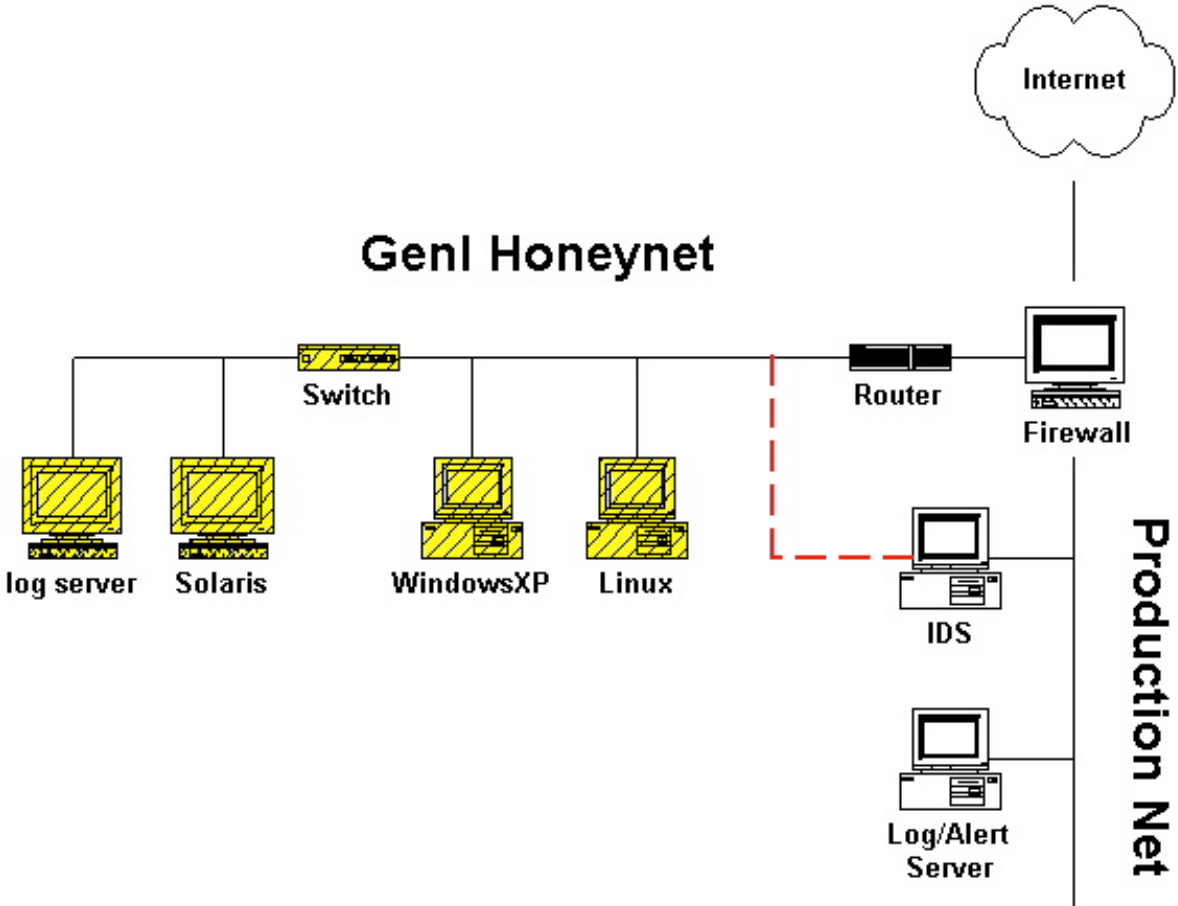
- A security resource whose value lies in being probed, attacked or compromised
- Anything going to or from a honeypot is likely a probe, attack or compromise

❖ Honeynet

- A network of honeypots
- All systems placed within the Honeynet are production systems : Solaris, Windows, Linux

<http://www.honeynet.org/alliance/>

Example of Honeynet



The threat is real!

- ❖ Computers scanned dozens of times a day
- ❖ Fastest time a honeypot manually compromised: 15 minutes (automatic, 92 seconds)
- ❖ Time before a default Linux Red Hat 6.2 successfully hacked is 72 hours
- ❖ 100% - 900% increase of activity from 2000 to 2001
- ❖ Its only getting worse

<http://www.honeynet.org/papers/stats>

Perspectives

❖ Data collection

- Several honeynets (different domains, locations, etc.)
- Need to analyze if data collected from different locations (e.g., .com vs. .edu) exhibit similar or different statistical patterns

❖ Data Analysis

- Identify attacks and characterize their distribution in space and time
 - known and new vulnerabilities
 - attack scenarios
 - trend analysis

❖ Security modeling and evaluation

- Take into account the lessons learnt from data
- Analyze how results are useful for designers/administrators

What we do NOT expect :

❖ Plausible attack rates / effort distribution

... necessary for "reliability / availability" measures

References

- ❖ L. Spitzner, *Honeypots: Tracking hackers*, 480p., ISBN 0321108957, Addison Wesley Professional, 2002. (<http://www.tracking-hackers.com/>)
- ❖ The Honeynet project (Ed.), *Know your enemy*, 352p., ISBN 0201746131, Addison-Wesley Pub Co, 2002. (<http://www.honeynet.org/>)
- ❖ R. Ortalo, Y. Deswarte, M. Kaâniche, "Experimenting with Quantitative Evaluation Tools for Monitoring Operational Security", *IEEE Transactions on Software Engineering*, Vol.25, N°5, pp.633-650, Sept./Oct. 1999.

